

QR Factorisation

Lemma: Given any two vectors $u, v \in \mathbb{R}^n$ with $\|u\|_2 = \|v\|_2 \exists w \in \mathbb{R}^n$ s.t. $H(w)u = v$.

Proof: take $w = r(u - v)$, any $r \in \mathbb{R} \setminus \{0\}$, so

$$\begin{aligned} w^T w &= r^2(u^T u - 2v^T u + v^T v) \\ &= 2r^2(u^T u - v^T u) && \text{as } \|u\| = \|v\| \\ &= 2r^2(u - v)^T u = 2rw^T u \end{aligned}$$

so $w^T u = (1/2r)w^T w$. Thus

$$\left(I - \frac{2}{w^T w} w w^T \right) u = u - \frac{2}{w^T w} \frac{w^T w}{2r} r(u - v) = v \quad \square$$

In particular

$$v = (\|u\|_2, 0, \dots, 0) = H(r[u_1 - \|u\|_2, u_2, \dots, u_n])u.$$

Can be applied to matrices:

if u is 1st column of A : write $H(u) = H_1, \alpha = \alpha_1$

$$H_1 A = \left[\begin{array}{c|ccc} \alpha_1 & \times & \dots & \times \\ \hline 0 & & & \\ \vdots & & & \\ 0 & & & \end{array} B \right] \quad \text{and if } H(\hat{w})B = \left[\begin{array}{c|ccc} \alpha_2 & \times & \dots & \times \\ \hline 0 & & & \\ \vdots & & & \\ 0 & & & \end{array} C \right]$$

then

$$H_2 = \left[\begin{array}{cc} 1 & 0 \\ 0 & H(\hat{w}) \end{array} \right] = H([0, \hat{w}])$$

satisfies

$$H_2 H_1 A = \left[\begin{array}{c|cc} \alpha_1 & \times & \times \dots \times \\ \hline 0 & \alpha_2 & \times \dots \times \\ \hline 0 & 0 & \\ \vdots & \vdots & \\ 0 & 0 & \end{array} C \right]$$

continuing inductively for n steps if $m > n$ gives

$$H_n \dots H_2 H_1 A = \begin{bmatrix} \alpha_1 & & * \\ & \ddots & \\ 0 & & \alpha_n \end{bmatrix} = R \in \mathbb{R}^{m \times n}$$

or for $m - 1$ steps if $m \leq n$ gives

$$H_{m-1} \dots H_2 H_1 A = \begin{bmatrix} \alpha_1 & & * \\ & \ddots & \\ 0 & & \alpha_m \end{bmatrix} = R \in \mathbb{R}^{m \times n}$$

Writing

$$Q = (H_n \dots H_2 H_1)^{-1} = H_1^T H_2^T \dots H_n^T = H_1 H_2 \dots H_n$$

as Householder matrices are symmetric gives

Theorem: Given any $A \in \mathbb{R}^{m \times n}$, \exists an orthogonal matrix $Q \in \mathbb{R}^{m \times m}$ and an upper triangular matrix $R \in \mathbb{R}^{m \times n}$ s.t.
 $A = QR$

Proof: Just take $H_i = I$ if i^{th} column is already zero below diagonal and the above procedure can not break down.

Remark: If $A = \begin{bmatrix} | & | & & | \\ a_1 & a_2 & \dots & a_n \\ | & | & & | \end{bmatrix}$

and $A = QR$

then $Q = \begin{bmatrix} | & | & & | \\ q_1 & q_2 & \dots & q_n \\ | & | & & | \end{bmatrix}$

and $\{q_1, q_2, \dots, q_n\}$ is an orthonormal basis for $\text{span}\{a_1, a_2, \dots, a_n\}$ if this set is linearly independent.

$\Rightarrow QR$ factorization essentially same as Gram-Schmidt.

Exercise: What happens to QR if $\{a_1, \dots, a_{n-1}\}$ is linearly independent with $a_n \in \text{span}\{a_1, \dots, a_{n-1}\}$?

Example: given data y_1, \dots, y_m at points x_1, \dots, x_m
find parameters in a model e.g. linear model $y = ax + b$
(parameters a, b)
such that

$\sum [y_i - (ax_i + b)]^2$ is min (regression)

same as

$$\min_{a, b} \left\| \begin{bmatrix} x_1 & 1 \\ x_2 & 1 \\ \vdots & \\ x_m & 1 \end{bmatrix} \begin{bmatrix} a \\ b \end{bmatrix} - \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_m \end{bmatrix} \right\|_2$$

QR Algorithm for eigenvalues (eig in matlab)

Set $A = A_1$

for $k = 1, 2, \dots$ $\left\{ \begin{array}{l} \text{factor } A_k = Q_k R_k \quad (\text{QR factorisation}) \\ \text{set } A_{k+1} = R_k Q_k \quad (\text{matrix multiply}) \end{array} \right.$

Lemma: $\{A_k\}$ are all similar matrices and so have same eigenvalues

Proof: $A_{k+1} = R_k Q_k = Q_k^T Q_k R_k Q_k = Q_k^{-1} (A_k) Q_k \quad \square$

Fact: $A_k \rightarrow$ upper triangular matrix as $k \rightarrow \infty$.

So for large k , $\text{diag}(A_k)$ are good approximations to the eigenvalues.

When complex conjugate eigenvalues arise, 2×2 real diagonal blocks remain in A_k which each have a pair of the complex conjugate eigenvalues.

Notes:

1. Speed of convergence depends on the size of gaps between the eigenvalues: more well separated \Rightarrow faster convergence.
2. Convergence is accelerated by the use of shifts (see problem sheet)
3. An orthogonal similarity reduction to Hessenberg form is always employed as a 1st step before apply the QR algorithm with shifts.

Reduction to Hessenberg form

Want a similarity transform $H_p \dots H_2 H_1 A H_1 H_2 \dots H_p$ where H_i are Householder matrices but if

$$H_1 A = \left[\begin{array}{c|ccc} \times & \times & \dots & \times \\ \hline \mathbf{0} & \times & \dots & \times \\ \vdots & \vdots & & \vdots \\ \mathbf{0} & \times & \dots & \times \end{array} \right]$$

then $H_1 A H_1$ is full i.e. postmultiplication destroys zeros created (otherwise a direct method for eigenvalues!)

So instead let

$$H_1 = \left[\begin{array}{c|ccc} 1 & 0 & \dots & 0 \\ \hline 0 & & & \\ \vdots & & & \\ 0 & & & \\ \hline & & & \end{array} \right]_{\substack{1 \\ n-1}} \quad , \quad A = \left[\begin{array}{c|c} a_1 & v_1^T \\ \hline u_1 & B_1 \end{array} \right] \text{ then}$$

$$H_1 A_1 H_1 =$$

$$\left[\begin{array}{cc} a_1 & v_1^T \\ K_1 u_1 & K_1 B_1 \end{array} \right] \left[\begin{array}{cc} 1 & 0 \\ 0 & K_1 \end{array} \right] = \left[\begin{array}{cc} a_1 & v_1^T K_1 \\ K_1 u_1 & K_1 B_1 K_1 \end{array} \right]$$

and choosing K_1 to be a Householder matrix satisfying $K_1 u_1 = (\alpha_1, 0, \dots, 0)^T$ we have

$$A^{(2)} = H_1 A H_1 = \left[\begin{array}{ccc} a_1 & b_1 & \hat{v}_2^T \\ \alpha_1 & a_2 & v_2^T \\ O & u_2 & B_2 \\ \hline 1 & 1 & n-2 \end{array} \right]_{\substack{1 \\ 1 \\ n-2}}$$

inductively (similar to before)

$$A^{(n-1)} = H_{n-2} \dots H_2 H_1 A H_1 H_2 \dots H_{n-2}$$

is in Hessenberg form

QR factorization can now be achieved by $n - 1$ Givens rotation matrices (see problem sheet) and the Hessenberg form is preserved by the QR algorithm ie. all of the A_k 's are upper Hessenberg.