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# Lecture 5: Newton's method for optimization problems (continued)

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C6.2/B2: Continuous Optimization

# Disadvantages of Newton's method for optimization

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- in the conditions of local convergence Theorem 9:  $x^k$  can get attracted to local maxima or saddle points of  $f$  if  $x^k$  sufficiently close to such points (as  $\nabla^2 f(x^*)$  only required to be nonsingular in Th 9).

**Example:**  $f : \mathbb{R} \rightarrow \mathbb{R}$ ,  $f(x) = -x^2$ ;

$x^* = 0$  is global maximizer;

apply Newton starting from  $x^0 = 1 \Rightarrow s^0 = -1$  ascent direction and  $x^1 = 0$ .

- Newton's method may fail to converge at all if  $x^0$  "too far" from solution (outside neighbourhood of local convergence, failure may occur).  
→ Newton is not globally convergent for general  $f$ .

# Disadvantages of Newton's method for optimization

Example of failure of Newton's method to converge globally.

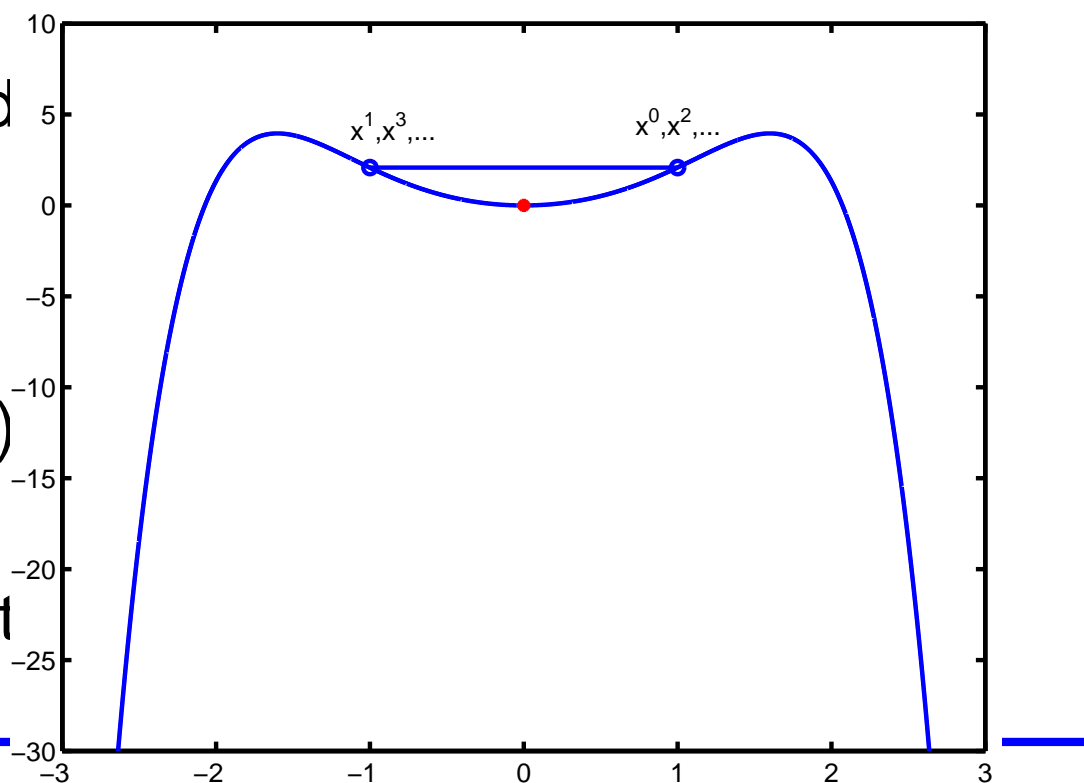
$$f : \mathbb{R} \rightarrow \mathbb{R}, \quad f(x) = -\frac{x^6}{6} + \frac{x^4}{4} + 2x^2.$$

$x^* = 0$  local minimizer;  $x = \pm\sqrt{(1 + \sqrt{17})/2} \approx \pm 1.6$  global max.

Newton's method applied to  $f$ , with  $x^0 = 1$ ;  
 $\Rightarrow x^{2k} = 1$  and  
 $x^{2k+1} = -1$ , for all  $k$ .

$-1$  and  $1$  are not (even) stationary points of  $f$ .

Note that  $s^k$  descent but we have gone "too far".



# Damped Newton's method

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⇒ include linesearch in Newton's method: **damped Newton**.

## Damped Newton's method for minimization:

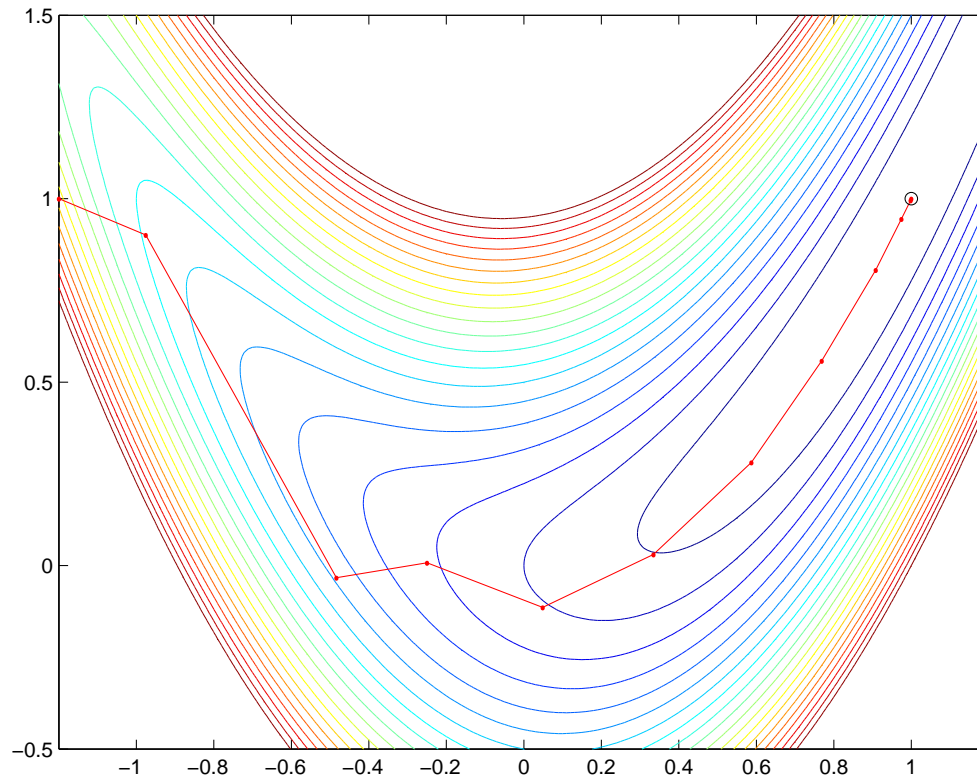
Choose  $\epsilon > 0$  and  $x^0 \in \mathbb{R}^n$ .

While  $\|\nabla f(x^k)\| > \epsilon$ , REPEAT:

- solve the linear system  $\nabla^2 f(x^k)s^k = -\nabla f(x^k)$ .
- set  $x^{k+1} = x^k + \alpha^k s^k$ , with  $\alpha^k \in (0, 1]$ ;  $k := k + 1$ . END.
- Damped Newton's method is a GLM provided  $\nabla^2 f(x^k)$  is positive definite so that  $s^k$  descent. Then  $\alpha^k$  can be computed by exact linesearch, bArmijo, etc.
- if  $\alpha^k \rightarrow 1$  as  $k \rightarrow \infty \implies$  damped Newton's method is locally quadratically convergent.
- (local convergence) Assume  $\nabla^2 f$  is Lipschitz cont., and  $\nabla^2 f(x^k) \succ 0$ . Let  $x^k \rightarrow x^*$  with  $\nabla^2 f(x^*) \succ 0$ . Let  $s^k =$ Newton direction in GLM and bArmijo linesearch have  $\beta < 0.5$  and  $\alpha_{(0)} = 1$ . Then,  $\alpha^k = 1$  for all  $k$  suff. large and  $x^k \rightarrow x^*$  quadratically.

# Local convergence for damped Newton with bArmijo

$$f(x_1, x_2) = 10(x_2 - x_1^2)^2 + (x_1 - 1)^2; \quad x^* = (1, 1).$$



Damped Newton with bArmijo linesearch applied to the Rosenbrock function  $f$ .

- $\beta < 0.5$  and  $\alpha_{(0)} = 1$  in bArmijo;  $\alpha^k = 1$  for suff. large  $k$ .

# Global convergence of damped Newton's method

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- recall backtracking Armijo (bArmijo) linesearch.

**Theorem 10** Let  $f \in \mathcal{C}^2(\mathbb{R}^n)$  be bounded below on  $\mathbb{R}^n$ .

Let  $\nabla f$  be Lipschitz continuous. Let the eigenvalues of  $\nabla^2 f(x^k)$  be positive and uniformly bounded below, away from zero (for all  $k$ ). Apply damped Newton's method to  $f$  with bArmijo linesearch and  $\epsilon = 0$ . Then

either

there exists  $l \geq 0$  such that  $\nabla f(x^l) = 0$

or

$\|\nabla f(x^k)\| \rightarrow 0$  as  $k \rightarrow \infty$ .  $\square$

- Theorem 10 is satisfied if  $f \in \mathcal{C}^2$  with  $\nabla f$  Lipschitz continuous is also strongly convex (i.e., the eigenvalues of  $\nabla^2 f(x)$  for all  $x$  are positive, bounded below, away from zero). Then  $s^k$  is descent for all  $k$ .
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# Global convergence of damped Newton's method ...

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**Proof of Theorem 10.** The conditions of Theorem 4 (Global convergence of GLM with bArmijo linesearch) are satisfied.

Thus Th 4 gives that either  $\exists l \geq 0$  such that  $\nabla f(x^l) = 0$  or

$$M_k := \min \left\{ \frac{|\nabla f(x^k)^T s^k|}{\|s^k\|}, |\nabla f(x^k)^T s^k| \right\} \longrightarrow 0 \text{ as } k \rightarrow \infty. \quad (\dagger)$$

Let  $\nabla^2 f(x^k) := H_k$ . Th assumptions on  $f \implies \forall s \in \mathbb{R}^n, s \neq 0,$

$$0 < \lambda_{\min} \leq \lambda_{\min}(H_k) \leq \frac{s^T H_k s}{\|s\|^2} \leq \lambda_{\max}(H_k) \leq \lambda_{\max}.$$

$$\begin{aligned} |\nabla f(x^k)^T s^k| &= |\nabla f(x^k)^T H_k^{-1} \nabla f(x^k)| \geq \lambda_{\min}(H_k^{-1}) \|\nabla f(x^k)\|^2 \\ &= \frac{\|\nabla f(x^k)\|^2}{\lambda_{\max}(H_k)} \geq \frac{\|\nabla f(x^k)\|^2}{\lambda_{\max}}. \end{aligned}$$

$$\|s^k\|^2 = \nabla f(x^k)^T H_k^{-2} \nabla f(x^k) \leq \lambda_{\max}(H_k^{-2}) \|\nabla f(x^k)\|^2 \leq \lambda_{\min}^{-2} \|\nabla f(x^k)\|^2.$$

$$\implies M_k \geq \min \left\{ \frac{\lambda_{\min}}{\lambda_{\max}} \|\nabla f(x^k)\|, \frac{1}{\lambda_{\max}} \|\nabla f(x^k)\|^2 \right\} \text{ for all } k$$

$$\implies \nabla f(x^k) \longrightarrow 0 \text{ as } k \rightarrow \infty. \quad \square$$

# Modified damped Newton methods

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If  $\nabla^2 f(x^k)$  is not positive definite, it is usual to solve instead

$$\left( \nabla^2 f(x^k) + M^k \right) s^k = -\nabla f(x^k),$$

where

- $M^k$  chosen such that  $\nabla^2 f(x^k) + M^k$  is “sufficiently” positive definite.
- $M^k := 0$  when  $\nabla^2 f(x^k)$  is “sufficiently” positive definite.

Options:

1. As  $\nabla^2 f(x^k)$  is symmetric, we can factor  $\nabla^2 f(x^k) = Q^k D^k (Q^k)^\top$ , where  $Q^k$  is orthogonal and  $D^k$  is diagonal, and set

$$\nabla^2 f(x^k) + M^k := Q^k \max(\epsilon I, |D^k|) (Q^k)^\top,$$

for some “small”  $\epsilon > 0$ . Expensive approach for large problems.

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# Modified damped Newton methods

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2. Estimate  $\lambda_{\min}(\nabla^2 f(x^k))$  and set

$$M^k := \max(0, \epsilon - \lambda_{\min}(\nabla^2 f(x^k)))I.$$

Cheaper. Often tried in practice but “biased” (may overemphasize a large negative eigval at the expense of small, positive ones).

3. Modified Cholesky: compute Cholesky factorization

$$\nabla^2 f(x^k) = L^k (L^k)^\top,$$

where  $L^k$  is lower triangular matrix. Modify the generated  $L^k$  if the factorization is in danger of failing (modify small or negative diagonal pivots, etc.).

Popular in computations.

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# Other directions for GLMs

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Choose/compute  $B^k$  to approximate  $\nabla^2 f(x^k)$ .

Let  $B^k$  **symmetric, positive definite** matrix. Let  $s^k$  be defined by

$$B^k s^k = -\nabla f(x^k).$$

Update  $B^k$  after the calculation of  $s^k$  and  $\alpha^k$ .

- $\implies s^k$  descent direction;
- $\implies s^k$  solves the problem

$$\text{minimize}_{s \in \mathbb{R}^n} m_k(s) = f(x^k) + \nabla f(x^k)^T s + \frac{1}{2} s^T B^k s.$$

- $s^k$  is a scaled steepest descent direction;
- Theorem 10 (global convergence) continues to hold with  $\nabla^2 f(x^k)$  replaced by  $B^k$  in the statement and proof.

# Approximating the Hessian matrix by finite differences

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Approximating the Hessian from gradient vals:  $i \in \{1, \dots, n\}$ ;

$$[\nabla^2 f(x)]e^i \approx \frac{1}{h}[\nabla f(x + he^i) - \nabla f(x)]$$

Cost of approximating  $\nabla^2 f(x)$  is  $n + 1$  gradient values.

For all finite-differencing, careful with the choice of  $h$  in computations:

- “too large”  $h \rightarrow$  inaccurate approximations,
- “too small”  $h \rightarrow$  numerical cancellation errors.

But successful techniques exist for smooth noiseless problems when sufficient function and/or gradient values can be computed.

For noisy problems, use **derivative-free optimization** methods (if problem size is not too large).

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# Quasi-Newton methods

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Secant approximations for computing  $B^k \approx \nabla^2 f(x^k)$

At the start of the GLM, choose  $B^0$  (say,  $B^0 := I$ ). After computing  $s^k = -(B^k)^{-1} \nabla f(x^k)$  and  $x^{k+1} = x^k + \alpha^k s^k$ , compute update  $B^{k+1}$  of  $B^k$ .

Wish list:

Compute  $B^{k+1}$  as a function of already-computed quantities  $\nabla f(x^{k+1}), \nabla f(x^k), \dots, \nabla f(x^0), B^k, s^k$ ,

$B^{k+1}$  should be symmetric, nonsingular (pos. def.),

$B^{k+1}$  “close” to  $B^k$ , a “cheap” update of  $B^k$ ,  $B^k \rightarrow \nabla^2 f(x^k)$ , etc.

$\implies$  a new class of methods: faster than steepest descent method, cheaper to compute per iteration than Newton's.

For the first wish, choose  $B^{k+1}$  to satisfy the secant equation

$$\gamma^k := \nabla f(x^{k+1}) - \nabla f(x^k) = B^{k+1}(x^{k+1} - x^k) = B^{k+1} \alpha^k s^k.$$

# Quasi-Newton methods ...

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## Interpretation of the secant equation:

It is satisfied by  $B^{k+1} := \nabla^2 f$  when  $f$  is a quadratic function.

The change in gradient contains information about the Hessian.

The gradient change predicted by the current quadratic model

$$\nabla f(x^{k+1}) - \nabla f(x^k) \approx \nabla q(x^k + \alpha^k s^k) - \nabla q(x^k) = -\alpha^k \nabla f(x^k),$$

where  $q(x^k + s) = f(x^k) + \nabla f(x^k)^\top s + \frac{1}{2} s^\top B^k s$

and  $s^k = -(B^k)^{-1} \nabla f(x^k)$ .

Want the new quadratic model

$$u(x^k + s) := f(x^k) + \nabla f(x^k)^\top s + \frac{1}{2} s^\top B^{k+1} s$$

to predict correctly the change in gradient  $\gamma^k$ , i.e.,

$$\gamma^k = \nabla f(x^{k+1}) - \nabla f(x^k) = \nabla u(x^{k+1}) - \nabla u(x^k) = B^{k+1} (x^{k+1} - x^k).$$

# Quasi-Newton methods ...

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Many ways to compute  $B^{k+1}$  to satisfy the secant equation.  
Trade-off between “wishes” on the list for some of the methods.

## Symmetric rank 1 updates.

[see Prob Sheet 3]

Set  $B^{k+1} := B^k + u^k (u^k)^\top$ , for some  $u^k \in \mathbb{R}^n$ , and all  $k \geq 0$ .

- $B^{k+1}$  symmetric, “close” to  $B^k$ .
- Work per iteration:  $\mathcal{O}(n^2)$  (as opposed to the  $\mathcal{O}(n^3)$  of Newton), due to Sherman-Morrison-Woodbury formula!

The secant equation  $\implies u^k = (\gamma^k - B^k \delta^k) / \rho^k$ ,  
where  $\delta^k := x^{k+1} - x^k = \alpha^k s^k$ ,  $(\rho^k)^2 := (\gamma^k - B^k \delta^k)^\top \delta^k > 0$ .

- $B^k$  may not be positive definite,  $s^k$  may not be descent.
- $\rho^k$  may be close to zero leading to large updates.

Other updates: **BFGS**, **DFP**, Broyden family, etc.

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# Quasi-Newton methods ...

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## BFGS updates.

[see Prob Sheet 3]

- Broyden-Fletcher-Goldfarb-Shanno (independently).

Set  $B_{k+1} := B_k + u_k u_k^\top + v_k v_k^\top$ , for some  $u_k \in \mathbb{R}^n$ ,  $v_k \in \mathbb{R}^n$ .

- It is a rank 2 update (if  $u_k$  and  $v_k$  are linearly independent).
- SWM formula yields  $\mathcal{O}(n^2)$  operations/iteration.
- In practice, update the Cholesky factors of  $B_k$  (still  $\mathcal{O}(n^2)$ ).

Given  $B_k = J_k J_k^\top$ , where  $J_k$  arbitrary nonsingular, and  $\|\cdot\|_F$  Frobenius norm, let  $J_{k+1}$  solve

$$\min_J \|J - J_k\|_F \quad \text{subject to} \quad J \delta_k = \gamma_k.$$

$$\Rightarrow B_{k+1} := J_{k+1} J_{k+1}^\top = B_k + u_k u_k^\top + v_k v_k^\top,$$

where  $u_k u_k^\top = -B_k \delta_k \delta_k^\top B_k / (\delta_k^\top B_k \delta_k)$ ,  $v_k v_k^\top = \gamma_k \gamma_k^\top / (\gamma_k^\top \delta_k)$ .

- Let  $J_k := L_k$  the lower triangular Cholesky factor of  $B_k$ .
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# Quasi-Newton methods ...

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## BFGS updates. (continued)

- Thus  $B_{k+1}$  is “close” to  $B_k$ .
  - $B_k$  symmetric pos. def.  $\Rightarrow B_{k+1}$  symmetric **pos. def.** (provided  $(\delta^k)^T \gamma^k > 0$ , ensured by say, **Wolfe linesearch**)
  - **BFGS method:** GLM with  $s_k := -B_k^{-1} \nabla f(x_k)$ , with  $B_k$  updated by BFGS formula on each iteration.
  - For global convergence of BFGS method, must use **Wolfe linesearch** to compute stepsize instead of bArmijo linesearch.
  - The BFGS method has **local Q-superlinear convergence!**
  - When applying the BFGS method with exact linesearches, to a strictly convex quadratic function  $f$ , then  $B_k = \nabla^2 f$  after  $n$  iterations.
  - Satisfies all the wishes on the wish list! Has been very popular when second derivatives of  $f$  are not available.
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# Appendix: providing derivatives to algorithms

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How to compute/provide derivatives to a solver?

- Calculate derivatives by hand when easy/simple objective and constraints; user provides code that computes them.
- Calculate or approximate derivatives automatically:
  - Automatic differentiation: breaks down computer code for evaluating  $f$  into elementary arithmetic operations + differentiate by chain rule. Software: ADIFOR, ADOL-C.
  - Symbolic differentiation: manipulate the algebraic expression of  $f$  (if available). Software: symbolic packages of MAPLE, MATHEMATICA, MATLAB.
  - Finite differencing  $\longrightarrow$  approximate derivatives.

See Nocedal & Wright, Numerical Optimization (2nd edition, 2006) for more details of the above procedures.