## Numerical Analysis

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Oxford Mathematical Institute

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 $0 = \det(A - \lambda I) = \det(P^{-1}(B - \lambda I)P) = \det(P^{-1})\det(P)\det(B - \lambda I),$ 

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so

$$det(A - \lambda I) = 0$$
 if, and only if,  $det(B - \lambda I) = 0$ .

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The basic **QR algorithm** is:

Set 
$$A_1 = A$$
.  
for  $k = 1, 2, \ldots$   
form the QR factorization  $A_k = Q_k R_k$   
and set  $A_{k+1} = R_k Q_k$   
end

## Proposition

The symmetric matrices  $A_1, A_2, \ldots, A_k, \ldots$  generated by the QR algorithm are all similar and thus have the same eigenvalues.

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The symmetric matrices  $A_1, A_2, \ldots, A_k, \ldots$  generated by the QR algorithm are all similar and thus have the same eigenvalues.

Proof. Since

$$A_{k+1} = R_k Q_k = (Q_k^{-1} Q_k) R_k Q_k = Q_k^{-1} (Q_k R_k) Q_k = Q_k^{-1} A_k Q_k,$$

 $A_{k+1}$  is similar to  $A_k$ . Furthermore, since  $Q_k^{-1} = Q_k^{\mathrm{T}}$ , the matrix  $A_{k+1}$  is symmetric if  $A_k$  is. By hypothesis,  $A_1 := A$  is symmetric; therefore each of the matrices  $A_k$ ,  $k = 1, 2, \ldots$ , is symmetric. This basic QR algorithm works since  $A_k \rightarrow$  a diagonal matrix as  $k \rightarrow \infty$ , the diagonal entries of which are the eigenvalues.

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However, a really practical, fast algorithm is based on some refinements.

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Note: direct reduction to triangular form would reveal the eigenvalues, but is not possible.

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Indeed, if

$$H(w)A = \begin{pmatrix} \times & \times & \cdots & \times \\ 0 & \times & \cdots & \times \\ \vdots & \vdots & \ddots & \vdots \\ 0 & \times & \cdots & \times \end{pmatrix}$$

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then  $H(w)AH(w)^{T}$  is generally full, i.e., all zeros created by pre-multiplication are destroyed by the post-multiplication.

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$$w = \left( \begin{array}{c} 0 \\ \hat{w} \end{array} \right) \ \, \text{where} \ \, H(\hat{w})u = \left( \begin{array}{c} \alpha \\ 0 \\ \vdots \\ 0 \end{array} \right),$$

it follows that

$$H(w)A = \begin{pmatrix} \gamma & u^{\mathrm{T}} \\ \alpha & \times & \vdots & \times \\ \vdots & \vdots & \vdots & \vdots \\ 0 & \times & \vdots & \times \end{pmatrix}, \quad ,$$

i.e., the  $\boldsymbol{u}^{\mathrm{T}}$  part of the first row of  $\boldsymbol{A}$  is unchanged.

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where 
$$B=H(\hat{w})CH^{\mathrm{T}}(\hat{w})$$
, as  $u^{\mathrm{T}}H(\hat{w})^{\mathrm{T}}=(lpha,~0,~\cdots,~0).$ 

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where  $B = H(\hat{w})CH^{\mathrm{T}}(\hat{w})$ , as  $u^{\mathrm{T}}H(\hat{w})^{\mathrm{T}} = (\alpha, 0, \cdots, 0)$ .

Note that  $H(w)AH(w)^{T}$  is symmetric as A is.

Now we inductively apply this to the smaller matrix B, as described for the QR factorization but using post- as well as pre-multiplications.

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The result of n-2 such Householder similarity transformations is the matrix

 $H(w_{n-2})\cdots H(w_2)H(w)AH(w)H(w_2)\cdots H(w_{n-2}),$ 

which is tridiagonal.

The QR factorization of a tridiagonal matrix can now easily be achieved with n-1 Givens rotations:

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The QR factorization of a tridiagonal matrix can now easily be achieved with n-1 Givens rotations: if A is tridiagonal

$$\underbrace{J(n-1,n)\cdots J(2,3)J(1,2)}_{Q^{\mathrm{T}}}A=R, \quad \text{upper triangular}.$$

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$$\underbrace{J(n-1,n)\cdots J(2,3)J(1,2)}_{Q^{\mathrm{T}}}A=R, \quad \text{upper triangular}.$$

Precisely, R has a diagonal and 2 super-diagonals,

$$R = \begin{pmatrix} \times & \times & \times & 0 & 0 & 0 & \cdots & 0 \\ 0 & \times & \times & \times & 0 & 0 & \cdots & 0 \\ 0 & 0 & \times & \times & \times & 0 & \cdots & 0 \\ \vdots & \vdots & & & & \vdots \\ 0 & 0 & 0 & 0 & \times & \times & \times & 0 \\ 0 & 0 & 0 & 0 & 0 & \times & \times & \times \\ 0 & 0 & 0 & 0 & 0 & 0 & \times & \times \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & \times \end{pmatrix}$$

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(exercise: check!).

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In the QR algorithm, the next matrix in the sequence is RQ.

In the QR algorithm applied to a tridiagonal matrix, the tridiagonal form is preserved when Givens rotations are used.

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Proof. If

$$A_k = QR = J(1,2)^{\mathrm{T}}J(2,3)^{\mathrm{T}}\cdots J(n-1,n)^{\mathrm{T}}R$$

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$$A_{k+1} = RQ = RJ(1,2)^{\mathrm{T}}J(2,3)^{\mathrm{T}}\cdots J(n-1,n)^{\mathrm{T}}.$$

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is tridiagonal, then

$$A_{k+1} = RQ = RJ(1,2)^{\mathrm{T}}J(2,3)^{\mathrm{T}}\cdots J(n-1,n)^{\mathrm{T}}.$$

Recall that:

post-multiplication of a matrix by  $J(i, i + 1)^{T}$  replaces columns i and i + 1 by linear combinations of the pair of columns, while leaving columns  $j \neq i, i + 1$  alone. Thus, since R is upper triangular, the only subdiagonal entry in  $RJ(1,2)^{T}$  is in position (2,1).

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Similarly, the only subdiagonal entries in

$$RJ(1,2)^{\mathrm{T}}J(2,3)^{\mathrm{T}} = (RJ(1,2)^{\mathrm{T}})J(2,3)^{\mathrm{T}}$$

are in positions (2,1) and (3,2).



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Inductively, the only subdiagonal entries in

$$RJ(1,2)^{\mathrm{T}}J(2,3)^{\mathrm{T}}\cdots J(i-2,i-1)^{\mathrm{T}}J(i-1,i)^{\mathrm{T}}$$
  
=  $(RJ(1,2)^{\mathrm{T}}J(2,3)^{\mathrm{T}}\cdots J(i-2,i-1)^{\mathrm{T}})J(i-1,i)^{\mathrm{T}}$ 

are in positions (j, j - 1),  $j = 2, \ldots i$ .

So, the lower triangular part of  $A_{k+1}$  only has nonzeros on its first subdiagonal.

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So, the lower triangular part of  $A_{k+1}$  only has nonzeros on its first subdiagonal.

However, then since  $A_{k+1}$  is symmetric, it must be tridiagonal.



## Using shifts.

One further and final step in making an efficient algorithm is the use of **shifts**:

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for 
$$k=1,2,\ldots$$
 form the QR factorization of  $A_k-\mu_k I=Q_k R_k$  and set  $A_{k+1}=R_k Q_k+\mu_k I$  end

For any chosen sequence of values of  $\mu_k \in \mathbb{R}$ ,  $\{A_k\}_{k=1}^{\infty}$  are symmetric and tridiagonal if  $A_1$  has these properties, and similar to  $A_1$ :



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$$\begin{aligned} A_2 &= R_1 Q_1 + \mu_1 I \\ &= Q_1^{\mathcal{T}} (Q_1 R_1) Q_1 + \mu_1 Q_1^{\mathcal{T}} Q_1 \\ &= Q_1^{\mathcal{T}} (A_1 - \mu_1 I) Q_1 + \mu_1 Q_1^{\mathcal{T}} Q_1 \\ &= Q_1^{\mathcal{T}} A_1 Q_1, \end{aligned}$$

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$$A_{3} = R_{2}Q_{2} + \mu_{2}I$$
  
=  $Q_{2}^{T}(Q_{2}R_{2})Q_{2} + \mu_{2}Q_{2}^{T}Q_{2}$   
=  $Q_{2}^{T}A_{2}Q_{2} = Q_{2}^{T}Q_{1}^{T}A_{1}Q_{1}Q_{2},$ 

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and similarly for  $A_4, A_5, \ldots$ 

The simplest shift to use is  $\mu_k = a_{n,n}$  for all k, which leads rapidly in almost all cases to

$$A_k = \left( \begin{array}{c|c} T_k & 0\\ \hline 0^T & \lambda \end{array} \right),$$

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where  $T_k$  is n-1 by n-1 and tridiagonal, and  $\lambda$  is an eigenvalue of  $A_1$ .

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Inductively, once this form has been found, the QR algorithm with shift  $a_{n-1,n-1}$  can be concentrated only on the  $(n-1) \times (n-1)$  leading submatrix  $T_k$ .

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This process is called **deflation**.

### The overall algorithm

for calculating the eigenvalues of an n by n symmetric matrix:

```
reduce A to tridiagonal form by orthogonal (Householder) similarity transformations.
```

for 
$$m = n, n - 1, \dots, 2$$
  
while  $a_{m-1,m} > \text{tol}$   
 $[Q, R] = \operatorname{qr}(A - a_{m,m} * I)$   
 $A = R * Q + a_{m,m} * I$   
end while  
record eigenvalue  $\lambda_m = a_{m,m}$   
 $A \leftarrow \text{leading } m - 1$  by  $m - 1$  submatrix of  $A$   
end

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record eigenvalue \lambda_1 = a_{1,1}
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