Definitions, Notation and Terminology

Notational Conventions

$[n] = \{1, 2, \ldots, n\}$.
$X^{(k)} = \{A \subseteq X : |A| = k\}$: set of k-element subsets of X.
For vertices $u \neq v$: $uv = \{u, v\} = vu$.
The *endvertices* or *ends* of an edge uv are u and v.

Graphs

A *graph* is an ordered pair (V, E) where $V \neq \emptyset$ is a finite (for now) set and $E \subseteq V^{(2)}$.

In a graph $G = (V, E)$:
the *vertex set* is $V = V(G)$ and *edge set* is $E = E(G)$,
the *order* $|G|$ of G is $|V|$ – the number of vertices, and
the *size* $e(G)$ of G is $|E|$ – the number of edges.
Vertices u, v are *adjacent* if $uv \in E$,
a vertex v and edge e are *incident* if v is an endvertex of e,
edges e and f *meet* if they share a vertex.
The *neighbourhood* of v is $\Gamma(v) = \Gamma_G(v) = \{u : uv \in E\}$ and,
the *degree* of v is $d(v) = d_G(v) = |\Gamma(v)|$.
v is *isolated* if $d(v) = 0$.

Isomorphism

An *isomorphism* from a graph G to a graph H is a bijection $\phi : V(G) \to V(H)$ such that $\phi(v)\phi(w) \in E(H)$ iff $vw \in E(G)$.
G and H are *isomorphic* if such a ϕ exists.

Subgraphs

A graph H is a *subgraph* of a graph G, written $H \subseteq G$, if $V(H) \subseteq V(G)$ and $E(H) \subseteq E(G)$.
If $W \subseteq V(G)$ then $G[W]$, the subgraph *induced* by W, is $(W, E(G) \cap W^{(2)})$, the graph formed by W and all edges of G with ends in W.
An *induced subgraph* of G is any such subgraph $G[W]$.
H is a *spanning subgraph* of G if $H \subseteq G$ and $V(H) = V(G)$.

Operations on graphs

The *complement* of $G = (V, E)$ is $\overline{G} = (V, V^{(2)} \setminus E)$.
A *non-edge* of G is an edge of $\overline{G}$.
For $e \in E$, the graph obtained by deleting e is $G - e = (V, E \setminus \{e\})$.
For $e \in V^{(2)} \setminus E$, the graph obtained by adding e is $G + e = (V, E \cup \{e\})$.
For $v \in V$, define $G - v = G[V \setminus \{v\}]$, i.e., delete v and any incident edges.

Standard graphs

K_n: complete graph on $n \geqslant 1$ vertices.

E_n: empty graph on $n \geqslant 1$ vertices.

P_n: path on $n \geqslant 0$ *edges* ($n+1$ vertices).

C_n: cycle on $n \geqslant 3$ vertices (also n edges).

$K_{a,b}$: complete bipartite graph with a vertices in one part and b in the other.

Formally:

$K_n = ([n], [n]^{(2)})$.

$E_n = ([n], \emptyset)$.

$P_n = (\{0, 1, \ldots, n\}, \{\{i-1, i\} : 1 \leqslant i \leqslant n\})$.

$C_n = ([n], \{12, 23, \ldots, \{n-1, n\}, n1\})$.

Further definitions

A graph G is *connected* if any two vertices are joined by a path/walk.

The *components* of G are the maximal connected subgraphs.

A *bridge* in G is an edge e whose deletion would disconnect the component of G containing e

A graph G is *bipartite* if we can partition the vertex set into $X \cup Y$ so that every edge is of the form xy, $x \in X$, $y \in Y$.

A graph is *acyclic* if it has no subgraph that is a cycle (i.e., is isomorphic to some C_n).

A *tree* is a connected acyclic graph.

A *forest* is an acyclic graph.

A *leaf* (in a tree/forest) is a vertex v with $d(v) = 1$.

If v is a vertex of $G = (V, E)$ and A and B are disjoint subsets of V we write

$\Gamma_A(v) = A \cap \Gamma(v)$ for the *neighbourhood of* v *in* A,

$d_A(v) = |\Gamma_A(v)|$ for the *degree of* v *into* A,

$e(A) = e(G[A])$ for the number of edges (of G) inside A and

$e(A, B)$ for the number of edges ab of G with $a \in A$ and $b \in B$.

Warnings!!!!

In some books P_n has n vertices, not n edges.

In some books 'graph' is used to mean 'multi-graph' – a variant where multiple edges between two vertices are allowed, and maybe edges from a vertex to itself. In most such books a 'simple graph' is what we call a graph.

Some people write $G \setminus e$ (*NOT* G/e) for $G - e$, and $G \setminus v$ for $G - v$.

You may also see $v(G)$ instead of $|G|$.

The term *size* is used in different ways by different people. Best to avoid and stick with $e(G)$ or 'number of edges'.

If you find an error please check the website, and if it has not already been corrected, e-mail riordan@maths.ox.ac.uk