

4. Composition series and the Jordan-Hölder theorem

Def. A sequence of subgroups

$$G_0 = \{e\} \triangleleft G_1 \triangleleft \dots \triangleleft G_{n-1} \triangleleft G_n = G$$

such that G_{i-1} is normal in G_i is called a composition series if each composition factor, G_{i+1}/G_i is simple.

Prop. Let G be a non-trivial and finite group.
Then G has a composition series.

proof. Proceed by induction on the size of the group, $|G|$.

If G is simple we are done: $G_0 = \{e\} \triangleleft G_1 = G$

Otherwise G has a non-trivial normal subgroup N with:

$|G/N| < |G|$ and $|N| < |G|$, and hence, by induction,

both have composition series:

$$\{e\} = Q_0 \triangleleft Q_1 \triangleleft \dots \triangleleft G/N = Q_k \quad \text{and} \quad 0 \triangleleft S_1 \triangleleft \dots \triangleleft S_{m-1} \triangleleft N = S_m$$

$$\text{Put } S_{i+m} = \pi^{-1}(Q_i)$$

Then

$$\{e\} \triangleleft S_1 \triangleleft \dots \triangleleft S_m \triangleleft S_{m+1} \triangleleft \dots \triangleleft G$$

is a composition series by the Third Isom. Th. $\square$

Ex: $C_{12} \triangleright C_6 \triangleright C_3 \triangleright \{e\}$
 $C_{12} \triangleright C_4 \triangleright C_2 \triangleright \{e\}$

$$C_{12}/C_6 \cong C_2, \quad C_6/C_3 \cong C_2, \quad C_3/\{e\} \cong C_3$$

$$C_{12}/C_4 \cong C_3, \quad C_4/C_2 \cong C_2, \quad C_2/\{e\} \cong C_2$$

$$S_4 \triangleright A_4 \triangleright V_4 \triangleright C_2 \triangleright \{e\}$$

Note: If G is finite abelian then all composition factors are cyclic of prime order.

Th (Jordan-Hölder theorem)

Let G be a finite, non-trivial group.

Then all composition series have the same length and, up to reordering, the same composition factors with multiplicities.

Proof: Assume we have composition series

$$G = G_r \triangleright \dots \triangleright G_0 = \{e\}$$

$$G = H_s \triangleright \dots \triangleright H_0 = \{e\}$$

Proceed by induction on r with $r \geq s$.

If $r=1$, $G = G_1$ is simple has only one comp. series.

Assume $r > 1$ and the theorem holds for any group having a comp series of length $r-1$

If $G_{r-1} = H_{s-1}$, then by assumption the comp. series for $G_{r-1} = H_{s-1}$ satisfy the theorem, and hence for $G_r = H_s$ as $G_r/G_{r-1} \cong H_s/H_{s-1}$ and $r=s$.

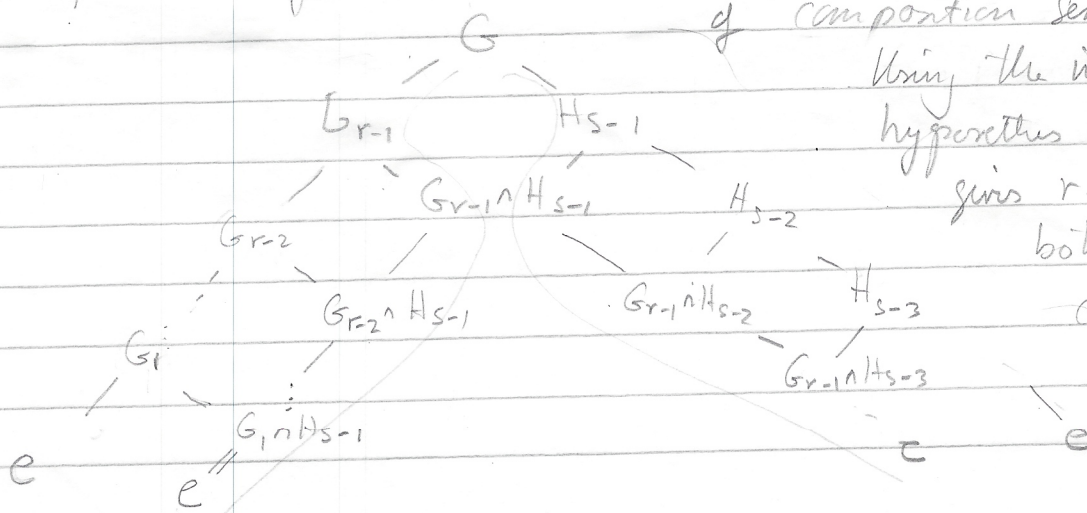
If $G_{r-1} \neq H_{s-1}$, consider

$$\begin{array}{ccccccc} G & \triangleright & G_{r-1} & \triangleright & G_{r-1} \cap H_{s-1} & \triangleright & \dots & \triangleright & G_0 = \{e\} \\ G & \triangleright & H_{s-1} & \triangleright & G_{r-1} \cap H_{s-1} & \triangleright & \dots & \triangleright & H_0 = \{e\} \end{array}$$

Note that $G = G_{r-1} H_{s-1}$ as

$$G_{r-1} H_{s-1} \triangleleft G_{r-1}$$

Repeated use of the Diamond Lemma gives an interlocking diagram of composition series.



Using the induction hypothesis on $G_{r-1} \cap H_{s-1}$ gives $r=s$ and both have the same composition factors.

5. Solvable groups

Def: For a group G and elements $x, y \in G$, define the commutator as

$$[x, y] = xyx^{-1}y^{-1}$$

and the derived group of G as

$$G' = \langle [x, y] \mid x, y \in G \rangle$$

Th: (1) G' is a normal subgroup of G . Furthermore, $G^{(n)}$ is normal in G .

(2) For $N \trianglelefteq G$, G/N is abelian iff $N \geq G'$

$G_{ab} = G/G'$ is the largest abelian quotient, the abelianization of G .

proof: (1) For any automorphism $\phi: G \rightarrow G$,

$$\phi(G') = G' \text{ as } \phi([x, y]) = \phi(x)\phi(y)\phi(x)^{-1}\phi(y)^{-1} = [\phi(x), \phi(y)]$$

G' is a
characteristic
subgroup of G

For $g \in G$, $c_g(x) := g^{-1}xg$ is an automorphism and hence $g^{-1}[x, y]g = [g^{-1}xg, g^{-1}yg] \in G'$.

(2) For $g_1, g_2 \in G$, $(g_1N)(g_2N) = (g_2N)(g_1N)$
iff $g_1g_2N = g_2g_1N$ iff $[g_1, g_2] \in N$

So G/N is abelian iff $N \geq G'$. □

Def: Let $G^{(1)} = G'$ and $G^{(n+1)} = (G^{(n)})'$. Then the derived series of G is

$$G \supseteq G' \supseteq G'' \supseteq \dots \supseteq G^{(n)} \supseteq \dots$$

G is solvable (or soluble) if $G^{(n)} = e$ for some n .

The first such n is called the derived length.

Prop: $G^{(n)}$ is normal in G for all $n \geq 1$.

proof: $G^{(n)} = (G^{(n-1)})'$ is characteristic in $G^{(n-1)}$.
By induction, $G^{(n-1)} \trianglelefteq G$ and c_g is an automorphism $G^{(n-1)} \rightarrow G^{(n-1)}$ for all $g \in G$.

Ex: $G = S_4$ with normal subgroups $A_4, V_4, \{e\}$;
 $G' = A_4$ as $S_4/A_4 \cong C_2$ is abelian but $S_4/V_4, S_4/\{e\}$ are not
 $G'' = (A_4)' = V_4$
 as $A_4/V_4 \cong C_3$ is abelian but $A_4/\{e\}$ is not
 $G^{(3)} = (V_4)' = \{e\}$
 as V_4 is abelian

So S_4 is solvable with derived series

$$S_4 \supseteq A_4 \supseteq V_4 \supseteq e \quad \text{of length 3}$$

Ex: A_5 is simple and not abelian. So $A_5' = A_5$.

so A_5 is not solvable: $S_5 \supseteq A_5 \supseteq A_5 \supseteq A_5 \dots$

Th: Let G be a finite group.

Then G is solvable if and only if every composition factor is abelian (i.e. cyclic of prime order).

proof: " $\Rightarrow$ " If G is solvable of length n then

$$(*) \quad G \supseteq G' \supseteq G'' \supseteq \dots \supseteq G^{(n-1)} \supseteq G^{(n)} = \{e\}$$

with $G^{(n-1)}/G^{(n)}$ abelian.

Using the
subgroup
correspondence

Each $G^{(k)}/G^{(n)}$ has composition series with composition factors all abelian, and $(*)$ can be refined to a composition series with all abelian factors.

" $\Leftarrow$ " Let $G = G_0 \supseteq G_1 \supseteq \dots \supseteq G_r = \{e\}$ will G_i/G_{i-1} be a composition series abelian.

Proceed by induction on k : to prove $G_k \supseteq G^{(k)}$ for $k=1, \dots, r$

G_0/G_1 is abelian, and so $G_1 \supseteq G'$ (by Th. (3))

Assume $G_k \supseteq G^{(k)}$. Then $G_k' \supseteq (G^{(k)})' = G^{(k+1)}$.

G_k/G_{k+1} is abelian, and so $G_{k+1} \supseteq G_k' \supseteq G^{(k+1)}$.

In particular, $G_r = \{e\} = G^{(r)}$ and G is solvable.

$G/G' = G_{ab}$ is the abelianization of G , i.e. the largest abelian quotient of G

$$\text{If } G = \langle X | R \rangle = F(X) / \langle\langle R \rangle\rangle$$

then

$$G_{ab} = \langle X | R \cup \{[x, y] \mid x, y \in X\} \rangle$$

$$\text{Ex: } \underline{B_{rn}} = \langle \sigma_1, \dots, \sigma_{n-1} \mid \begin{array}{l} \sigma_i \sigma_j = \sigma_j \sigma_i \quad |i-j| = 1 \\ \sigma_i \sigma_{i+1} \sigma_i = \sigma_{i+1} \sigma_i \sigma_{i+1} \end{array} \rangle$$

$(B_{rn})_{ab}$: add relations $\sigma_i \sigma_j = \sigma_j \sigma_i$ for all i, j ;
hence, the braid relation gives
 $\sigma_i = \sigma_{i+1}$ for all i

$$\Rightarrow (B_{rn})_{ab} \cong \langle \sigma_1 \rangle \cong \mathbb{Z}$$

$$\text{and } B_{rn}' = \ker \left(\phi : B_{rn} \rightarrow \mathbb{Z} \right) \\ \sigma_i \mapsto 1$$

$$\text{Ex: } \underline{F_n} = \langle a_1, \dots, a_n \mid \emptyset \rangle \quad \text{free group on } n\text{-generators}$$

$$(F_n)_{ab} = \langle a_1, \dots, a_n \mid [a_i, a_j] \rangle \cong \mathbb{Z}^n$$

Ex₃: If G is abelian, then $G = G_{ab}$
and its derived length is 1.

6. Semi-direct Products

Def: Let G be a group, $H \leq G$, $N \trianglelefteq G$.

Then G is an internal semi-direct product of H and N if $G = NH$ and $N \cap H = \{e\}$.

Write

$$G = N \rtimes H$$

(equivalently, $G = H \ltimes N$)

Note: In this case $G/N \cong H$.

Ex: Direct-product $G = H \times N$.

Ex: $D_{2n} = C_n \rtimes C_2$

$$S_n = A_n \rtimes C_2$$

$$S_4 = V_4 \rtimes S_3$$

Recall: $n_1, n_2 \in N$, $h_1, h_2 \in H$

$$(n_1, h_1)(n_2, h_2) = n_1(h_1 n_2 h_1^{-1})h_1 h_2$$

↑ multiplication is twisted by conjugation by h_1 .

$$= n_1 c_{h_1}(n_2) h_1 h_2$$

$$c_h(x) = h x h^{-1} \quad c_h \in \text{Aut}(G)$$

Lemma: $N \trianglelefteq G$, $H \leq G$

Then $H \rightarrow \text{Aut}(N)$ via $h \mapsto c_h$ is a homom.

proof: We have already seen that

$$G \rightarrow \text{Aut}(G) \quad \text{is a homomorphism.}$$

$$g \mapsto c_g$$

Hence, the restriction to $H \leq G$ is a homomorphism.

As N is normal in G , for all $g \in G$ (hence all $h \in H$)

$$c_g(n) = g n g^{-1} \in N \quad \text{for all } n \in N.$$

So c_g restricts to an element in $\text{Aut}(N)$.

Def/Th: Given two groups H and N , and a homomorphism $\phi: H \rightarrow \text{Aut}(N)$

define a group, the semi-direct product $N \rtimes_{\phi} H$ with elements (n, h) $n \in N, h \in H$ and product $(n_1, h_1)(n_2, h_2) = (n_1 \phi(h_1)(n_2), h_1 h_2)$.
 Note $N \times \{e\} \trianglelefteq N \rtimes_{\phi} H$ and $\{e\} \times H \leq N \rtimes_{\phi} H$.

Ex: $N = C_n$
 $\text{Aut}(C_n) \cong C_n^{\times}$ group of units in $C_n \cong \mathbb{Z}/n$
 $= \{k \mid (n, k) = 1\}$

$n=3$: $\text{Aut}(C_3) \cong \{1, 2\} \cong C_2$

$n=12$: $\text{Aut}(C_{12}) \cong \{1, 5, 7, 11\} \cong C_2 \times C_2$

Ex: $N = C_3 = \langle x \rangle$ $H = C_4 = \langle y \rangle$

How many semi-direct products $C_3 \rtimes C_4$

$\phi: C_4 \rightarrow \text{Aut}(C_3) = C_2 = \{1, \tau\}$, $\tau(x) = x^2 = x^{-1}$

ϕ is determined by $\phi(y)$

$\phi(y) = 1$ then $C_3 \rtimes_{\phi} C_4 = C_3 \times C_4$

$\phi(y) = \tau$ then $C_3 \rtimes C_4 = \langle x, y \mid x^3 = y^4 = e, yx = x^{-1}y \rangle$

Prop: H is normal in $G = N \rtimes_{\phi} H$ iff ϕ is trivial i.e. $G = N \times H$

proof: $n^{-1}hn \in H \Leftrightarrow n^{-1}\phi(h)(n)h \in H \Leftrightarrow n^{-1}\phi(h)(n) \in H$
 $\Leftrightarrow \phi(h)(n) = n$ as $N \cap H = \{e\}$

N
SI

Q
SI

(11)

$$K \hookrightarrow G \twoheadrightarrow G/K$$

short exact sequence

Def: Let N and Q be two groups.

An extension of N by Q is a group G such that G has a normal subgroup K such that

$$K \cong N \text{ and } G/K \cong Q.$$

The extension is split if there is a subgroup $H \cong Q$ of G such that $H \hookrightarrow G \twoheadrightarrow G/K \cong Q$ is an isomorphism.

Prop: G is a split extension of N by Q iff there exists $\phi: Q \rightarrow \text{Aut}(N)$ such that

$$G \cong N \rtimes_{\phi} Q.$$

Ex: C_6 is an extension of C_2 by C_3 ; it is split.
 C_8 is an extension of C_2 by C_4 ; it is not split.
 S_n is an extension of A_n by C_2 ; it is split.

Proof: If $G = N \rtimes_{\phi} Q$ then $N \hookrightarrow G \xrightarrow{\pi} Q$ is split with $s: Q \rightarrow G$ the inclusion of the subgroup.
 If $N \hookrightarrow G \twoheadrightarrow Q$ is split, we may identify Q via s with a subgroup of G .

Then $\phi: Q \rightarrow \text{Aut}(N)$ via $q \mapsto c_q: N \rightarrow N$
 $m \mapsto q n q^{-1}$

defines a homomorphism and $N \rtimes_{\phi} Q$ is well-defined.

Consider $N \rtimes_{\phi} Q \xrightarrow{\alpha} G$ via $\alpha(n, q) = nq$.

• α is a homomorphism:

$$\begin{aligned} \alpha((n_1, q_1)(n_2, q_2)) &= \alpha(n_1 q_1 n_2 q_1^{-1}, q_1 q_2) = n_1 q_1 n_2 q_1^{-1} q_1 q_2 \\ &= n_1 q_1 n_2 q_2 = \alpha(n_1, q_1) \alpha(n_2, q_2) \end{aligned}$$

• α is injective: $\alpha(n, q) = e \iff n = q^{-1} \in N \cap Q$
 $\iff n = q^{-1} = e$ as s is a splitting and π is injective on $s(Q)$.

• α is surjective: every $g \in G$ is contained in some $g \ker \pi$; but $G/\ker \pi \cong Q \subseteq G$ are the coset representatives.