

**Definition 1.** A relation  $R$  on a class  $A$  is well-founded if and only if

$$\forall x \subseteq A [x \neq \emptyset \rightarrow \exists m \in x \forall b \in x \neg [bRm]]$$

(i.e. every non-empty subset of  $A$  has an  $R$ -minimal element).

**Definition 2.** A relation  $R$  on a class  $A$  is set-like if and only if

$$\forall x \in A \exists z \forall t [t \in z \leftrightarrow tRx]$$

**Lemma 1.** Suppose  $R$  is a set-like relation on a class  $A$ . Then  $\text{pred}_R: A \rightarrow U$  given by  $\text{pred}_R(x) = \{t \in A : tRx\}$  is a class function on  $A$ .

*Proof.* By **Extensionality** and the definition of being set-like.  $\square$

**Lemma 2.** Suppose  $R$  is a well-founded, set-like relation on a class  $A$ . Then  $R^*$  given by

$$xR^*y \equiv \exists n \in \omega \exists f: n+2 \rightarrow A [f(0) = x \wedge f(n+1) = y \wedge \forall m \in n+1 [f(m)Rf(m+1)]]$$

is a well-founded, set-like relation on  $A$  that is transitive and contains  $R$ .

*Proof.* Clearly it is a relation on  $A$ .

If  $xRy$  then  $n = 0$  and  $f = \{\langle 0, x \rangle, \langle 1, y \rangle\}$  witnesses  $xR^*y$  so  $R \subseteq R^*$ .

For  $x, y \in A$  such that  $xR^*y$ , we write  $p(x, y)$  for the smallest  $n \in \omega$  which witnesses  $xR^*y$ .

We will show transitivity (i.e.  $\forall x, y, z \in A [xR^*y \wedge yR^*z \rightarrow xR^*z]$ ) by induction on  $p(y, z)$ . So assume this is true provided  $p(y, z) < n$  and that  $p(y, z) = n$ . Let  $g$  be a witness of  $yR^*z$  with  $\text{dom}(g) = n+2$ . The special case  $n = 0$  is easy to handle manually. Otherwise, note that  $p(y, g(n)) < n$  so that by inductive hypothesis there is  $m \in \omega$  and  $f: m+2 \rightarrow A$  witnessing  $xR^*g(n)$ . Then  $m+1$  together with  $f \cup \{\langle m+2, z \rangle\}$  witnesses  $xR^*z$ .

Next we show that  $R^*$  is set-like. Fix  $x \in A$ . By induction on  $n \in \omega$ , we can see that

$$P_n = \{z \in A : zR^*x \wedge p(z, x) \leq n+1\} = \bigcup \{\text{pred}_R(y) : yR^*x \wedge p(y, x) \leq n\}$$

is a set (by **Union** and **Replacement**) and (another application of **Union** and **Replacement**) so is  $\text{pred}_{R^*}(x) = \bigcup \{P_n : n \in \omega\}$ .

Finally for well-foundedness, let  $C \subseteq A$  and  $C \neq \emptyset$  and assume that  $x$  has no  $R^*$ -minimal element. Let  $x_0 \in C$  and

$$a = \{x \in \text{pred}_{R^*}(x_0) \cup \{x_0\} : \exists y \in C yR^*x\}$$

(noting that this is a set by **Separation** and the fact that  $R^*$  is set-like.)

Since  $x_0 \in C$  is not  $R^*$ -minimal,  $x_0 \in a$ , so  $a \neq \emptyset$ . Since  $R$  is well-founded,  $a$  has an  $R$ -minimal element  $m$ . By definition of  $a$ , we can find  $y \in C$  such that  $yR^*m$ . Since  $mR^*x_0$  (or  $m = x_0$ ) and  $R^*$  is transitive, we obtain  $yR^*x_0$  and since  $C$  has no  $R^*$ -minimal element there must be  $y' \in C$  with  $y'R^*y$ , giving  $y \in a$ . Find an  $n \in \omega$  and  $f: n+2 \rightarrow A$  witnessing  $yR^*m$ . Now  $f(n)R^*m$  and  $mR^*x_0$  or  $m = x_0$  given  $f(n)R^*x_0$  by transitivity of  $R^*$ . Also  $yR^*f(n)$  or  $y = f(n)$  so that  $f(n) \in a$ . But we also have  $f(n)Rm$ , contradicting  $R$ -minimality of  $m$  in  $a$ .  $\square$

**Definition 3.** Suppose  $R$  is a well-founded, set-like relation on  $U$  and  $F: U \times U \rightarrow U$  is a class function.

We define

$$\begin{aligned} \psi_{F,R}(a, g) &\equiv g \text{ is a function on } \{a\} \cup \text{pred}^*(a) \wedge \\ &\quad \forall x \in \text{dom}(g) \ g(x) = F(x, g|_{\text{pred}(x)}) \end{aligned}$$

and

$$G_{F,R} \equiv \{\langle a, b \rangle : \exists g \ \psi_{F,R}(a, g) \wedge b = g(a)\}$$

**Theorem 1.** Suppose  $R$  is a well-founded, set-like relation on  $U$ .

If  $F: U \times U \rightarrow U$  is a class function then, writing  $G$  for  $G_{F,R}$ ,  $G$  is a class function on  $U$  such that

$$\forall x \ G(x) = F(x, G|_{\text{pred}(x)})$$

and for every class function  $H$  on  $U$  such that  $\forall x \ H(x) = F(x, H|_{\text{pred}(x)})$  we have  $\forall x \ G(x) = H(x)$ .

We first prove a lemma (this is sort of the ‘internal’ version of the theorem ‘up to any  $a$ ’).

**Lemma 3.** Suppose  $R$  is a well-founded, set-like relation on  $Y$ .

If  $F: U \times U \rightarrow U$  is a class function then

$$\forall a \ \exists! g \ \psi_{F,R}(a, g).$$

*Proof.* Suppose this is not the case. Fix  $x_0 \in U$  such that  $\neg \exists! g \ \psi_{F,R}(x_0, g)$  and let  $a$  be  $R^*$ -minimal in  $\{y \in \{x_0\} \cup \text{pred}^*(x_0) : \neg \exists! g \ \psi_{F,R}(y, g)\}$  (noting that this is a set and contains  $x_0$  so is non-empty). Note that  $a$  is then  $R^*$ -minimal such that  $\neg \exists! g \ \psi_{F,R}(a, g)$ .

**Case  $\neg \exists g \ \psi_{F,R}(a, g)$ :** Note that by  $R^*$ -minimality of  $a$ , we have

$$\forall y \in \text{pred}^*(a) \ \exists! g \ \psi_{F,R}(y, g)$$

so that

$$f = \{g : \exists y \in \text{pred}^*(a) \ \psi_{F,R}(y, g)\}, \quad h = \bigcup f$$

are sets (by **Replacement** and **Union**). Note that if  $g, g' \in f$  and  $y \in \text{dom}(g) \cap \text{dom}(g')$  then  $\hat{y} = \{y\} \cup \text{pred}^*(y) \subseteq \text{dom}(g) \cap \text{dom}(g')$  and that  $\psi_{F,R}(y, g|_{\hat{y}})$  as well as  $\psi_{F,R}(y, g'|_{\hat{y}})$ . Further  $y R^* a$  so that by  $R^*$ -minimality of  $a$  we must have  $g|_{\hat{y}} = g'|_{\hat{y}}$  and hence  $g(y) = g'(y)$ . Thus  $h$  is a function.

Next for  $y \in \text{pred}^*(a)$  find  $g$  such that  $\psi_{F,R}(y, g)$  so that  $g \subseteq h$  and hence  $y \in \text{dom}(g) \subseteq \text{dom}(h)$ .

Thus  $h$  is a function on  $\text{pred}^*(a)$  and we note that

$$g = h \cup \{\langle a, F(a, h|_{\text{pred}(a)}) \rangle\}$$

is a function on  $\{a\} \cup \text{pred}^*(a)$  such that  $\psi_{F,R}(a, g)$ , a contradiction.

**Case  $\exists g, g'$  [ $\psi_{F,R}(a, g) \wedge \psi_{F,R}(a, g') \wedge g \neq g'$ ]:** By  $R^*$ -minimality of  $a$ ,  $\forall y \in \text{pred}^*(a)$   $g(y) = g'(y)$  and in particular  $g|_{\text{pred}(a)} = g'|_{\text{pred}(a)}$ . But then  $g(a) = F(a, g|_{\text{pred}(a)}) = F(a, g'|_{\text{pred}(a)}) = g'(a)$  so that  $g = g'$ , a contradiction.  $\square$

*Proof of the General Recursion Theorem.* First, if  $x \in U$  suppose that there are  $y, y' \in U$  such that  $\langle x, y \rangle, \langle x, y' \rangle \in G_{F,R}$ . Find  $g, g'$  such that  $\psi_{F,R}(x, g)$  and  $g(x) = y$  and  $\psi_{F,R}(x, g')$  and  $g'(x) = y'$ . Then by the Lemma  $g = g'$  so that  $y = g(x) = g'(x) = y'$ . Thus  $G_{F,R}$  is a class function on some class  $A$ .

Next if  $x \in U$  then there is (by the Lemma)  $g$  such that  $\psi_{F,R}(x, g)$  so that  $\langle x, g(x) \rangle \in G_{F,R}$ . Hence  $G_{F,R}$  is a class function on  $U$ .

That  $G(a) = F(a, G|_{\text{pred}(a)})$  follows from the construction of  $G$  and the method in the previous Lemma.

Similarly the ‘uniqueness’ is an easy consequence of the previous Lemma by considering a ‘ $R^*$ -minimal’ counterexample.  $\square$