Numerical Analysis Sheet 2 — HT21 SVD, least-squares and eigenvalues

form the same matrix $\begin{bmatrix} c & s \\ -s & c \end{bmatrix}$). Show that if $x \in \Re^n$ then (denoting by J(i, j) a Givens rotation)

$$J(1,n)J(1,n-1)\cdots J(1,3)J(1,2) x = \begin{bmatrix} \beta \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

for some β and further prove that $\beta^2 = x^T x$.

2. Consider the least-squares problem

$$\min_{x} \|Ax - b\|, \qquad A \in \mathbb{R}^{m \times n}, b \in \mathbb{R}^{m \times 1}, m \ge n.$$
(1)

Suppose that $\operatorname{rank}(A) = n$.

- (a) Derive a solution x based on the (full or thin) QR factorisation of A.
- (b) Show that $x = (A^T A)^{-1} A^T b$ is also the solution for (1). (The first QR-based method has better numerical stability and hence preferred)
- 3. Show that all the eigenvalues of a real symmetric matrix A are real, and eigenvalues of an orthogonal matrix are on the unit circle.
- 4. What is the SVD of a normal matrix A, with respect to the eigenvalues and eigenvectors? What if A is symmetric? And unitary?
- 5. If $A \in \mathbb{R}^{n \times n}$ is nonsingular, what is the SVD of A^{-1} in terms of that of A?

6. Give estimates based on Gershgorin's theorem for the eigenvalues of

$$A = \begin{bmatrix} 9 & 1 & 0 \\ 1 & 4 & \varepsilon \\ 0 & \varepsilon & 1 \end{bmatrix}, \quad |\varepsilon| < 1.$$

Find a way to establish the tighter bound $|\lambda_3 - 1| \leq \varepsilon^2$ on the smallest eigenvalue of A. (Hint: consider diagonal similarity transformations.)

- 7. We have seen how the QR algorithm computes the eigenvalue decomposition of symmetric matrices. Using QR, describe an algorithm that computes the SVD of $A \in \mathbb{R}^{m \times n} (m \ge n)$. You can assume that $\operatorname{rank}(A) = n$.
- 8. The (unshifted) QR algorithm performs: $A_1 = A$, and for k = 1, 2, ... form the QR factorisation $A_k = Q_k R_k$, and set $A_{k+1} = R_k Q_k$.

Let A be a 10×10 symmetric matrix with eigenvalues 1, 2, ..., 10. Which (off-diagonal) elements of A_k converge to 0 the fastest? (Recall the connection between the QR algorithm and power method).

- 9. (Computational) Implement the QR algorithm in MATLAB or Python, and explore the following. (this exercise should be doable in about 10 lines of code!)
 - (a) Verify the above problem in your code. (take $A = Q\Lambda Q^T$, where Q is a randomly generated orthogonal matrix, e.g. [Q,R]=qr(randn(n));).
 - (b) What will happen if the QR algorithm (without shifts) is applied to an orthogonal matrix A? Explain why this is not a 'counterexample' for the convergence of the QR algorithm. Modify the algorithm so that it computes an eigenvalue decomposition of A.
- 10. (Optional) Let $A \in \mathbb{C}^{n \times n}$ and define

$$C_{i,j} = \left\{ z \in \mathbb{C} \left| |a_{ii} - z| |a_{jj} - z| \le \left(\sum_{\substack{k \neq i \\ k=1}}^n |a_{ik}| \right) \left(\sum_{\substack{k \neq j \\ k=1}}^n |a_{jk}| \right) \right\}, \quad 1 \le i < j \le n.$$

These are called the *ovals of Cassini*.

- (a) Prove that all eigenvalues of A lie in the union of the ${}_{n}C_{2} = \frac{n(n-1)}{2}$ regions $\cup_{i < j}C_{i,j}$.
- (b) Prove that $C := \bigcup_{i < j} C_{i,j} \subseteq D$, that is, C is a subset of the union of the Gerschgorin disks $D := \bigcup_i \{ z \in \mathbb{C} | |z a_{ii}| \le \sum_{j \neq i} |a_{ij}| \}.$