

## 15. Another example: Algebraically closed fields

**15.1** Let  $\mathcal{L}_r = \{+, \times, 0, 1\}$  be the language of rings. In this language one can formulate axioms that say that a given  $\mathcal{L}_r$  structure is a ring, or that it is a field:

$$\psi_1 : \forall x \forall y \, x + y = y + x$$

$$\psi_2 : \forall x \forall y \forall z \, (x + y) + z = x + (y + z)$$

$$\psi_3 : \text{etc}$$

### 15.2 There are many examples of fields:

$\langle A; +, \times, 0, 1 \rangle$  where  $A$  is:  $\mathbb{Q}$ ,  $\mathbb{R}$ ,  $\mathbb{C}$ , and also **finite fields**  $\mathbb{F}_p = \mathbb{Z}/p\mathbb{Z}$ , where  $p$  is a prime number, and their extensions (see later).

These are not elementarily equivalent: e.g.  $\mathbb{C}$  has a solution to  $x^2 = -1$  but  $\mathbb{R}$  doesn't. Both solve  $x^2 = 2$  but  $\mathbb{Q}$  doesn't, while a finite field has  $\mathbb{F}_p$  **characteristic**  $p$ , i.e.  $p = 0$  while  $\mathbb{Q}$ ,  $\mathbb{R}$ ,  $\mathbb{C}$  have characteristic 0 (Definition: no positive integer  $n = 0$ ).

**15.3** A field  $K$  contains a **prime field**, smallest subfield, which is either  $\mathbb{F}_p$  if  $\text{char}(K) = p > 0$  or  $\mathbb{Q}$  if  $\text{char}(K) = 0$ .

**15.4** One can **extend** fields algebraically, like  $\mathbb{Q}(\sqrt{2})$  extends  $\mathbb{Q}$ , or  $\mathbb{C} = \mathbb{R}(i)$  extends  $\mathbb{R}$ , or by **transcendental elements**, like the field of rational functions  $\mathbb{Q}(x)$  extends  $\mathbb{Q}$ .

**15.5** If  $K \subset L$  are fields, a **transcendence basis** for  $L$  over  $K$  is a set  $X \subset L$  of algebraically independent elements such that  $L$  is algebraic over  $K(X)$ . Any two of these have the same **cardinality** (that's a theorem), called the **transcendence degree** of the extension  $L/K$ .

### **15.6. Examples.**

$$\text{tr.deg.}(\mathbb{Q}(\sqrt{2})/\mathbb{Q}) = 0,$$

$$\text{tr.deg.}(\mathbb{C}/\mathbb{R}) = 0,$$

$$\text{tr.deg.}(\mathbb{Q}(x)/\mathbb{Q}) = 1,$$

$$\text{tr.deg.}(\mathbb{C}/\mathbb{Q}) = |\mathbb{C}| = 2^{\aleph_0}.$$

**15.7 Theorem (Steinitz).** *Two algebraically closed fields are isomorphic if they have same prime field and same transcendence degree over the prime field.*

**“Proof”.** By back and forth construction of an isomorphism starting with transcendence bases  $X_1, X_2$  of each field  $K_1, K_2$  over prime field  $F$ .  $F(X_1), F(X_2)$  are isomorphic under any bijection  $X_1 \rightarrow X_2$ . Further elements are algebraic. Taking an element of  $K_1$  (say) we can solve the “same” equation in  $K_2$ .

Thus we see that any algebraically closed field of characteristic 0 and transcendence degree over  $\mathbb{Q}$  equal to  $2^{\aleph_0}$  is **isomorphic** to  $\mathbb{C}$ .

**15.8.** We can axiomatize the theory of algebraically closed fields of characteristic  $p \geq 0$  (recall 10.9). Characteristic 0 is axiomatized by the sequence of formulas

$$1 + \dots + 1 \neq 0 \text{ (with } n \text{ 1's, } n = 2, 3, \dots)$$

Algebraic closedness can be axiomatized by a sequence of axioms (for each degree  $d = 2, \dots$ ) saying that every monic polynomial of degree  $d$  has a root:

$$\forall a_0 \dots \forall a_{d-1} \exists x (a_0 + a_1 x + \dots + a_{d-1} x^{d-1} + x^d = 0)$$

We get **Theories**  $\text{ACF}_0$ ,  $\text{ACF}_p$ .

**15.9.** By (a stronger version of) Löwenheim-Skolem,  $\text{ACF}_0$  is **complete** (max consistent) and is the theory of  $\text{Th}(\langle \mathbb{C}, +, \times, 0, 1 \rangle)$ . Likewise  $\text{ACF}_p$  is complete.

This means that a theorem about  $\mathbb{C}$  that can be expressed in  $\mathcal{L}_r$  holds in any algebraically closed field of char 0, and has a proof in FOPC, even if it was proved using methods of analysis.

**15.10.** We can also prove things about  $\mathbb{C}$  by going to positive characteristic! (i.e.  $p > 0$ )

Any proof in  $\text{ACF}_0$  can only use finitely many of the axioms  $1 + \dots + 1 \neq 0$  and hence goes through in any algebraically closed field whose characteristic is sufficiently large.

**15.11. Theorem.** *Let  $\phi$  be a sentence in  $\mathcal{L}_r$ .  
TFAE:*

- \*  $\mathbb{C} \models \phi$
- \*  $\phi$  holds in **all** alg closed fields of char 0
- \* ... **some** ...
- \* For arbitrarily large  $p$ ,  $\phi$  holds in some  $\text{ACF}_p$
- \*  $\phi$  holds in any  $\text{ACF}_p$  provided  $p > N(\phi)$

One rather surprising application:

**15.12. Theorem** (Ax-Grothendieck) *Consider a polynomial map  $F : \mathbb{C}^n \rightarrow \mathbb{C}^n$ . If  $F$  is injective then it is surjective.*

**“Proof”** A finite field  $\mathbb{F}_p$  has a unique extension of each finite degree  $n$ , denoted  $\mathbb{F}_q, q = p^n$ . Likewise  $\mathbb{F}_q, q = p^n$ , and the algebraic closure  $\overline{\mathbb{F}}$  of  $\mathbb{F}_p$  is the union of these. If we consider a polynomial map  $F : \overline{\mathbb{F}}^n \rightarrow \overline{\mathbb{F}}^n$  then the coefficients will be in  $\mathbb{F}_q$  for some (large enough)  $q$ , and then  $F$  maps  $\mathbb{F}_q^n \rightarrow \mathbb{F}_q^n$  for this field and all its extensions. These fields are **finite**, so if a map is injective it must be surjective (and *vice versa*). So the statement holds in positive characteristic. But it is expressible in  $\mathcal{L}_r$ : for each  $n, d$ , with  $F$  a “general” degree  $d$  poly map:

$\forall \text{ coefficients}(F \text{ injective} \rightarrow F \text{ surjective})$

So it holds in (is a theorem of)  $\text{ACF}_0$ .

## 16. Normal Forms

### (a) Prenex Normal Form

A formula is in **prenex normal form (PNF)** if it has the form

$$Q_1x_{i_1}Q_2x_{i_2}\cdots Q_rx_{i_r}\psi,$$

where each  $Q_i$  is a quantifier (i.e. either  $\forall$  or  $\exists$ ), and where  $\psi$  is a formula containing no quantifiers.

#### 16.1 PNF-Theorem

*Every  $\phi \in \text{Form}(\mathcal{L})$  is logically equivalent to an  $\mathcal{L}$ -formula in **PNF**.*

*Proof:* Induction on  $\phi$   
(working in the language with  $\forall, \exists, \neg, \wedge$ ):

$\phi$  atomic: OK

$$\phi = \neg\psi,$$

$$\text{say } \phi \leftrightarrow \neg Q_1 x_{i_1} Q_2 x_{i_2} \cdots Q_r x_{i_r} \chi$$

$$\text{Then } \phi \leftrightarrow Q_1^- x_{i_1} Q_2^- x_{i_2} \cdots Q_r^- x_{i_r} \neg\chi,$$

where  $Q^- = \exists$  if  $Q = \forall$ , and  $Q^- = \forall$  if  $Q = \exists$

$$\phi = (\chi \wedge \rho) \text{ with } \chi, \rho \text{ in PNF}$$

$$\text{Note that } \vdash (\forall x_j \psi[x_j/x_i] \leftrightarrow \forall x_i \psi),$$

provided  $x_j$  does not occur in  $\psi$  (Ex. 12.5)

So w.l.o.g. the variables quantified over in  $\chi$  do not occur in  $\rho$  and vice versa.

$$\text{But then, e.g. } (\forall x \alpha \wedge \exists y \beta) \leftrightarrow \forall x \exists y (\alpha \wedge \beta) \text{ etc.}$$

□

## (b) Skolem Normal Form

**Recall:** In the proof of CT, we introduced witnessing new constants for existential formulas such that

$\exists x \phi(x)$  is satisfiable iff  $\phi(c)$  is satisfiable.

This way an  $\exists x$  in front of a formula could be removed at the expense of a new constant.

Now we remove existential quantifiers ‘inside’ a formula at the expense of extra function symbols:

### 16.2 Observation:

*Let  $\phi = \phi(x, y)$  be an  $\mathcal{L}$ -formula with  $x, y \in \text{Free}(\phi)$ . Let  $f$  be a new unary function symbol (not in  $\mathcal{L}$ ).*

Then  $\forall x \exists y \phi(x, y)$  is satisfiable iff  $\forall x \phi(x, f(x))$  is satisfiable.

( $f$  is called a **Skolem function** for  $\phi$ .)

*Proof:* ' $\Leftarrow$ ': clear

' $\Rightarrow$ ': Let  $\mathcal{A}$  be an  $\mathcal{L}$ -structure with  $\mathcal{A} \models \forall x \exists y \phi(x, y)$

$\Rightarrow$  for every  $a \in A$  there is some  $b \in A$  with  $\phi(a, b)$

Interpret  $f$  by a function assigning to each  $a \in A$  one such  $b$

(this uses the Axiom of Choice!).  $\square$

**16.3. Example:**  $\mathbb{R} \models \forall x \exists y (x \dot{=} y^2 \vee x \dot{=} -y^2)$ .

Here  $f(x) = \sqrt{|x|}$  will do.

## 16.4. Theorem

*For every  $\mathcal{L}$ -formula  $\phi$   
there is a formula  $\phi^*$   
(with new constant and function symbols)  
having only universal quantifiers in its PNF  
such that*

*$\phi$  is satisfiable iff  $\phi^*$  is.*

*More precisely,  
any  $\mathcal{L}$ -structure  $\mathcal{A}$   
can be made into a structure  $\mathcal{A}^*$   
interpreting the new constant and function sym-  
bols  
such that*

*$\mathcal{A} \models \phi$  iff  $\mathcal{A}^* \models \phi^*$ .*