

This is a preliminary problem sheet, to get the ball rolling. There is a ‘bonus’ problem for MFoCS students overleaf. Hints/solutions will be put on the website near the start of week 2. Problem sheet 1 (based on the first two week’s lectures) will be for the first class, in week 3 or 4.

*Estimates and asymptotics, union bound and first-moment method*

1. Prove the following inequalities:

- (a)  $1 + x \leq e^x$  for all real  $x$ .
- (b)  $(1 + a)^n \leq e^{an}$  for  $a > -1$ ,  $n \geq 0$ .
- (c)  $k! \geq k^k / e^k$  for  $k \geq 1$ .
- (d)  $\left(\frac{n}{k}\right)^k \leq \binom{n}{k} \leq \frac{n^k}{k!} \leq \left(\frac{en}{k}\right)^k$  for  $1 \leq k \leq n$ .

2. For the following functions  $f(n)$  and  $g(n)$ , decide whether  $f = o(g)$  or  $g = o(f)$  or  $f = \Theta(g)$  as  $n \rightarrow \infty$ :

- (a)  $f(n) = \binom{n}{k}$ ,  $g(n) = n^k$ , first for  $k$  fixed and then for the case where  $k = k(n) \rightarrow \infty$  as  $n \rightarrow \infty$ ;
- (b)  $f(n) = (\log n)^{1000}$ ,  $g(n) = n^{1/1000}$ ;

3. In lectures we saw that the  $k$ th diagonal Ramsey number satisfies

$$R(k, k) > n - \binom{n}{k} 2^{1 - \binom{k}{2}},$$

for each integer  $n$ . By considering  $n = \lfloor e^{-1} k 2^{k/2} \rfloor$ , deduce that

$$R(k, k) \geq (1 - o(1)) e^{-1} k 2^{k/2}.$$

4. Show that if  $n, k, \ell \geq 1$  are integers and  $0 < p < 1$ , then

$$R(k, \ell) > n - \binom{n}{k} p^{\binom{k}{2}} - \binom{n}{\ell} (1 - p)^{\binom{\ell}{2}}.$$

5. Let  $H$  be an  $r$ -uniform hypergraph with fewer than  $\frac{3^{r-1}}{2^r}$  edges. Prove that the vertices of  $H$  can be coloured using three colours in such a way that in each edge, all three colours are represented.
6. Let  $F$  be a collection of binary strings (“codewords”) of finite length, where the  $i$ th codeword has length  $c_i$ . Suppose that no member of  $F$  is a prefix of another member (so you can decode any string made up by concatenating codewords as you go along, without looking ahead). Show that  $\sum_i 2^{-c_i} \leq 1$  (the *Kraft inequality* for prefix-free codes).

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*Bonus question (for MFoCS students, optional for others):*

A (finite, or infinite and convergent) sum  $S = \sum_{i \geq 0} a_i$  is said to *satisfy the alternating inequalities* if the partial sum  $\sum_{i=0}^t a_i$  is at least  $S$  for all even  $t$  and at most  $S$  for all odd  $t$ ; that is, the partial sums alternately over- and under-estimate the final result.

7. Let  $I_1, \dots, I_n$  be the indicator functions of  $n$  events  $E_1, \dots, E_n$ . For  $0 \leq r \leq n$  let  $S_r = \sum_{A \subseteq [n], |A|=r} \prod_{i \in A} I_i$ , where  $[n] = \{1, 2, \dots, n\}$ . Show that

$$\prod_{i=1}^n (1 - I_i) = \sum_{r=0}^n (-1)^r S_r, \quad (0.1)$$

and that the sum satisfies the alternating inequalities. [ Both sides are random; the statement is that the relevant inequalities *always* hold. You may want to consider different cases according to how many of the events  $E_i$  hold. ] Deduce that

$$\mathbb{P}(\text{no } E_i \text{ holds}) = \sum_{r=0}^n (-1)^r \sum_{A \subseteq [n], |A|=r} \mathbb{P}(\cap_{i \in A} E_i), \quad (0.2)$$

and that the sum satisfies the alternating inequalities. [This is a form of the inclusion-exclusion formula.]