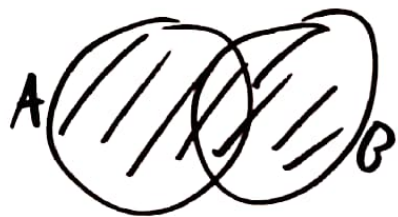


§1 The first moment method.

The union bound

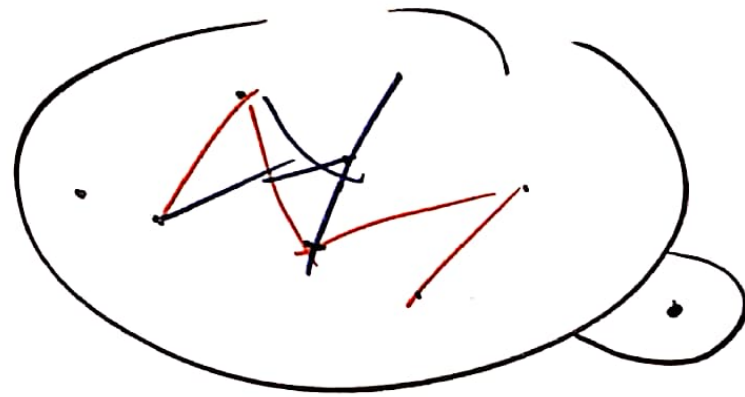
$$IP(A \cup B) \leq IP(A) + IP(B)$$



If $A_1, \dots, A_m$ are events

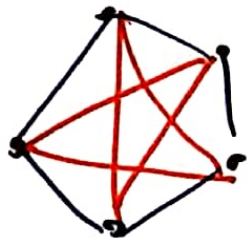
$$IP(\cup_{i=1}^m A_i) \leq \sum_{i=1}^m IP(A_i).$$

For $k, l \geq 1$ integers the Ramsey number $R(k, l) = \min \{n : \text{every red/blue colouring of the edges of } K_n \text{ contains a red } K_k \text{ or a blue } K_l\}$



x x x x x x ✓ ✓ ✓ ✓ ✓ ✓
n n+1

$x \ x \ x \ \dots \ x \ \checkmark \ \checkmark \ \checkmark$
 n



$$R(3,3) > 5$$

Theorem 1.1 (Erdős 1947).

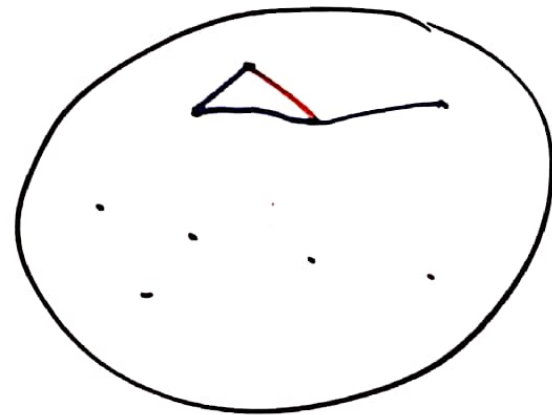
Let $n, k \geq 1$ be integers s.t.

$$\binom{n}{k} 2^{1 - \binom{k}{2}} < 1. \quad \text{Then } R(k,k) > n.$$

Pt Suffices to show K_n has
 a 'good' colouring, i.e. a

colouring with no
 monochromatic K_k .

Colour edges of K_n randomly
 each red with prob. $\frac{1}{2}$, blue
 otherwise. with the colours of
 the edges independent.



Let A_i be the event
that i^{th} copy of K_k is
monochromatic, $i = 1, 2, \dots, \binom{n}{k}$

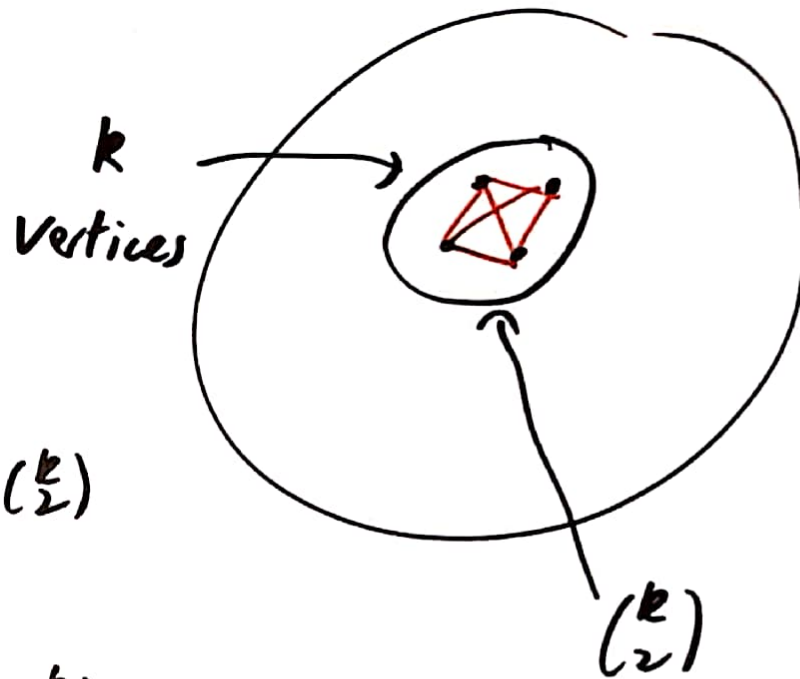
$$IP(A_i) = \left(\frac{1}{2}\right)^{\binom{k}{2}} + \left(\frac{1}{2}\right)^{\binom{k}{2}} = 2^{1-\binom{k}{2}}$$

$$IP(\cup A_i) \leq \sum_{i=1}^n IP(A_i) = \binom{n}{k} 2^{1-\binom{k}{2}} < 1$$

ie. $IP(\text{random colouring is 'bad'}) < 1$

ie. $IP(\text{random colouring is 'good'}) > 0$

$\therefore \exists$ a good colouring $\square$



Theorem 1.1 If $n, k \geq 1$ are integers s.t.

$$\binom{n}{k} 2^{1-\binom{k}{2}} < 1 \text{ then } R(k, k) > n.$$

Cor 1.2 For $k \geq 3$, $R(k, k) \geq 2^{k/2}$.

Pf Let $n = \lfloor 2^{k/2} \rfloor$. Then

$$\binom{n}{k} 2^{1-\binom{k}{2}} \leq \frac{(2^{k/2})^k}{k!} 2^{1-\binom{k}{2}} = \frac{2^{k^2/2}}{k!} 2^{1-\frac{k^2}{2}+\frac{k}{2}} = \left(\frac{2^{1+\frac{k}{2}}}{k!} \right)^{\times \sqrt{2}} < 1 \text{ for } k \geq 3$$

$\times (k+1)$

$$\therefore R(k, k) > \lfloor 2^{k/2} \rfloor \quad \square$$

decreases as $k \nearrow$ ($k \geq 2$)

$$\binom{n}{k} = \frac{n!}{k! (n-k)!} = \frac{n(n-1) \dots (n-k+1)}{k!} \leq \frac{n^k}{k!}$$

$$\binom{k}{2} = \frac{k(k-1)}{2} = \frac{k^2}{2} - \frac{k}{2}$$

§1 The first moment method

If X is a RV (random variable) its first moment is

just its mean / ~~E~~ expectation $E[X] = \sum_x x P(X=x) = E(X) = EX$

key property: E is LINEAR

If X and Y are RVs, $E[X+Y] = \overset{\textcircled{X}}{E[X]} + E[Y]$

If λ is a constant $E[\lambda X] = \lambda E[X]$.

$\textcircled{X}$ ALWAYS HOLDS - no independence or other condition.

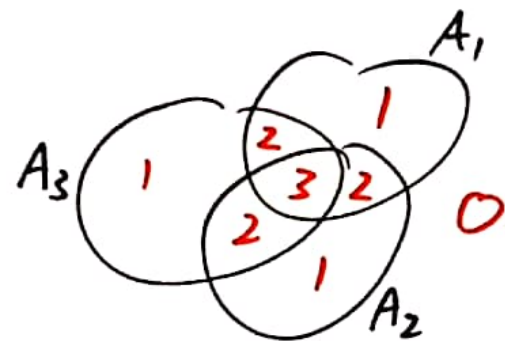
If A is an event, its indicator function I_A or 1_A is the RV taking value 1 when A holds, 0 otherwise.

$$E[I_A] = 1 \cdot P(A) + 0 \cdot P(A^c) = P(A).$$

Let $A_1, \dots, A_n$ be events and let $I_1, \dots, I_n$ be their indicator functions. ($I_j = I_{A_j}$)

Let $X = \sum_{j=1}^n I_j$ counts how many of the A_j hold.

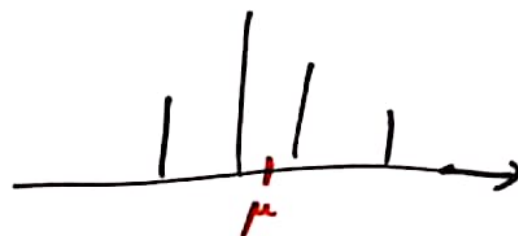
$$\text{Then } E[X] = \sum_{j=1}^n E[I_j] = \sum_{j=1}^n P(A_j).$$



Observation: if X is a RV with $E[X] = \mu$

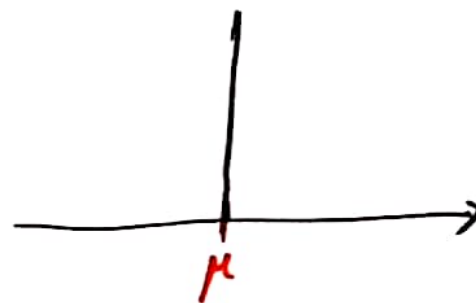
then $IP(X \geq \mu) > 0$

Similarly $IP(X \leq \mu) > 0$.



Theorem 1.3 Let $n, k \geq 1$ be integers.

Then $R(k, k) > n - \binom{n}{k} 2^{1 - \binom{k}{2}}$.

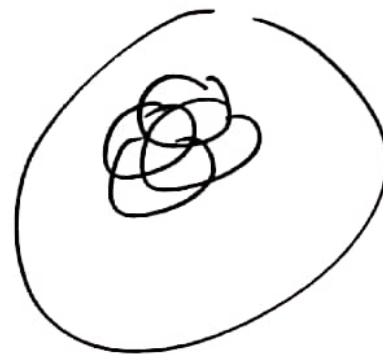


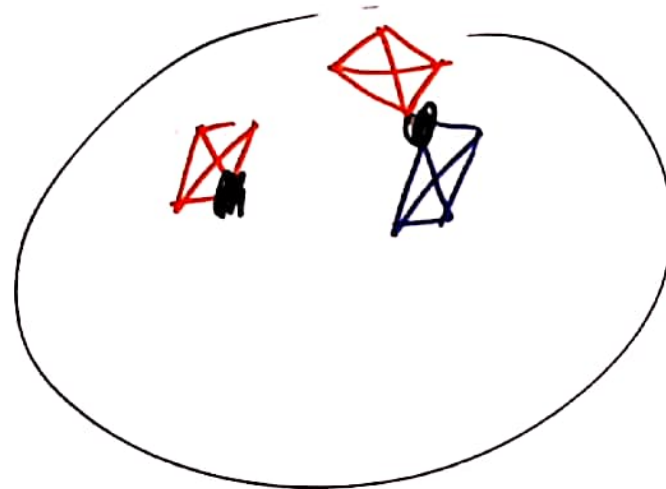
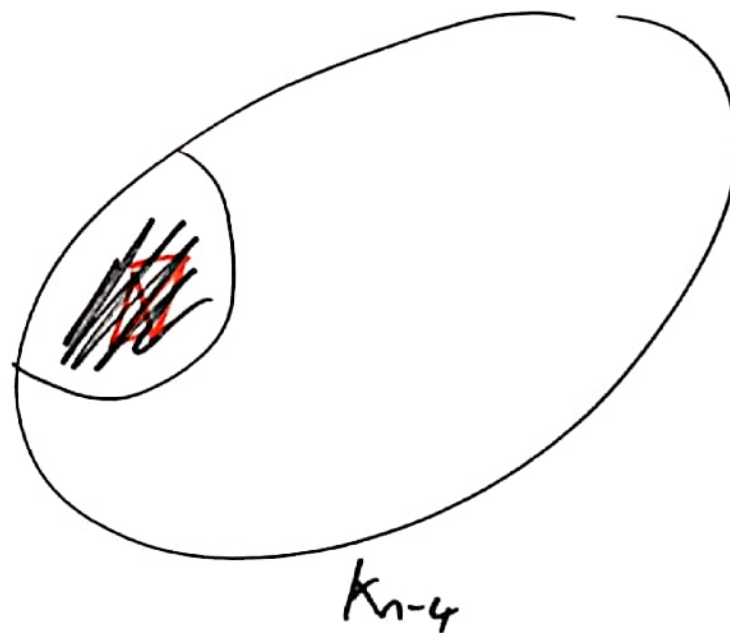
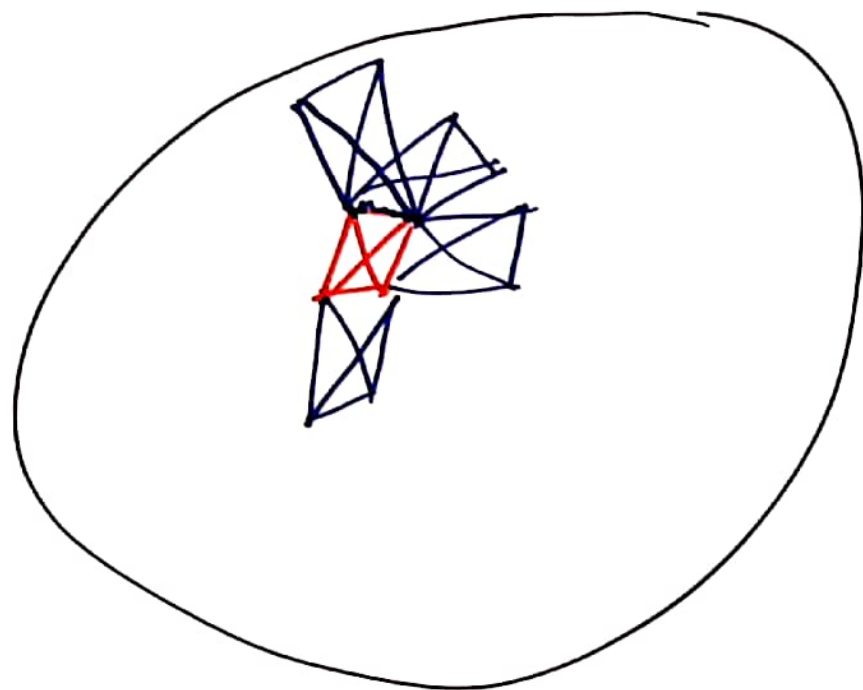
Pf Colour edges of K_n red/blue as before.

Let $X = \#$ monochromatic copies of K_k

Define A_i as before

Then $E[X] = \sum IP(A_i) = \binom{n}{k} 2^{1 - \binom{k}{2}} =: \mu$





By observation $\mathbb{P}(X \leq \mu) > 0$

ie. a colouring with $t \leq \mu$ defects (mono. K_k)

Pick such a colouring Pick a vertex of each defect
to delete.

In total delete $\leq t \leq \mu$ vertices

Obtain a 'good' colouring of K_m where $m \geq n - \mu$

$$\therefore R(k, k) > m$$

□

Theorem 1.3 Let $k, n \geq 1$ be integers

Then $R(k, k) > n - \binom{n}{k} 2^{1 - \binom{k}{2}}$.

Cor 1.4 $R(k, k) \geq (1 - o(1)) \frac{k}{e} 2^{k/2}$ as $k \rightarrow \infty$.

i.e. $\forall \epsilon > 0 \exists k_0 \forall k \geq k_0 \quad R(k, k) \geq (1 - \epsilon) \frac{k}{e} 2^{k/2}$

Pf Exercise. Take $n = \left\lfloor \frac{k}{e} 2^{k/2} \right\rfloor$ $\square$

Remarks on calculation

Considering structures of a certain type.

Aim: show $\exists$ a good one.

$$\left. \begin{array}{l} \text{Total \# structures} \geq \text{some formula} \\ \text{\# bad structures} \leq \text{some other formula} \end{array} \right\} \Rightarrow \exists \text{ good structure.}$$

$\vee$

A set $S \subseteq \mathbb{R}$ or $\mathbb{Z}$ is sum-free if $\nexists a, b, c \in S$

s.t. $a + b = c$. eg. $\{1, 2, 6\}$ is NOT sum-free

as $1 + 1 = 2$

$\{1, 3, 7, 9, 25\}$ is sum-free.

Consider $[n] = \{1, 2, \dots, n\}$. How big a sum-free set
can we find in $[n]$?

Around $\frac{n}{2}$ — odd numbers
 \ large numbers

$\underbrace{\quad}$
 $\odot \cdot \odot \cdot \odot \cdot \odot \mid \dots \cdot \cdot$
 $1 \quad 3 \quad 5 \quad 7 \quad \frac{n}{2}$

S a set of integers: can we find large
sum-free $A \subseteq S$?

eg try $A = \{a \in S : a \text{ odd}\}$

. $\frac{m}{2}$ m

Theorem 1.5 (Erdős 1965). Let S be a set of $n \geq 1$ non-zero integers.

There is some $A \subseteq S$ s.t. A is sum-free and $|A| > \frac{n}{3}$.

Pf Let $S = \{s_1, \dots, s_n\}$.

Trick: work mod p prime, p of the form $3k+2$

choose $p > 2 \max |S|$ so $s_1, \dots, s_n$ are distinct & $\neq 0$
mod p .

A s.t. $A \subseteq \mathbb{Z}$ is sum-free mod p if

$$\nexists a, b, c \in A \text{ s.t. } a + b \equiv c \pmod{p}$$

$$\Updownarrow \\ a + b = c$$

$$\text{SUM-FREE mod } p \Rightarrow \text{SUM-FREE}$$

For $r=1, 2, \dots, p-1$ the map
 $a \mapsto ra \pmod p$ is a bijection
 from $\mathbb{Z}_p^* = \{1, 2, \dots, p-1\}$ to itself.

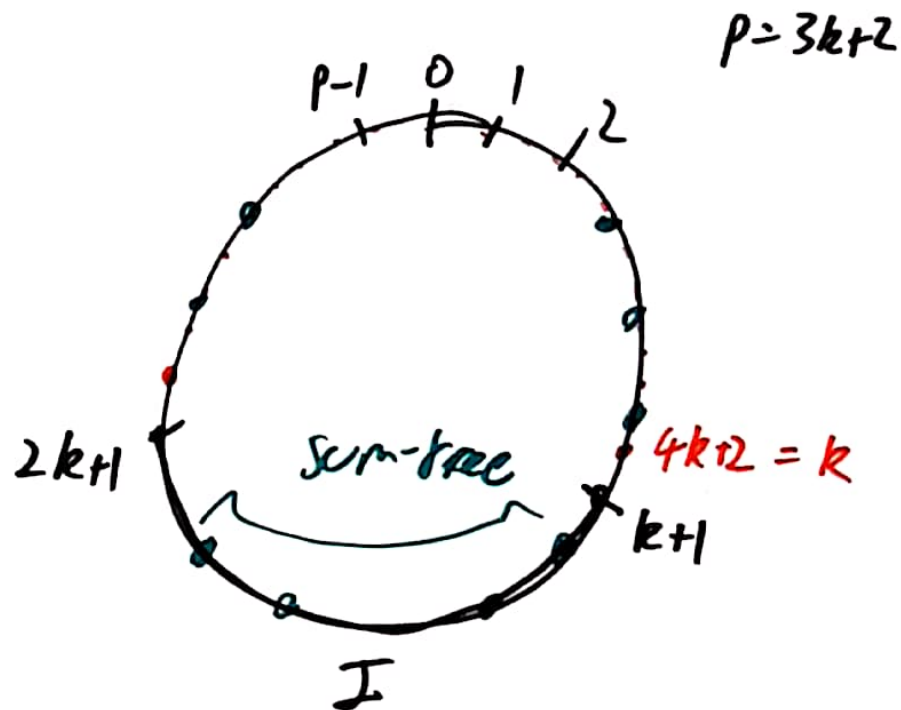
Let R be random with

$$\mathbb{P}(R=r) = \frac{1}{p-1} \quad r=1, 2, \dots, p-1$$

$$\text{Let } X = \# \{i : R_{Si_n}^{\pmod p} \in \mathcal{I}\}$$

$$\mathbb{E}[X] = \sum_{i=1}^n \mathbb{P}(R_{Si_n}^{\pmod p} \in \mathcal{I})$$

Fix i : Since $x \mapsto x s_i$ is a bijection $\mathbb{Z}_p^*$ to itself



$\mathcal{I}$ is sum-free mod p .

Rs_i takes all values in $\mathbb{Z}_p^*$ with equal probability.

$$\therefore \mathbb{P}(Rs_i \bmod p \in \mathcal{I}) = \frac{|\mathcal{I}|}{p-1} = \frac{k+1}{3k+1} > \frac{1}{3}$$

$$\therefore \mathbb{E}[X] = \sum = n \times \text{[circled]} > \frac{n}{3}$$

$\therefore \exists$ value r s.t. when $R=r$, $X > \frac{n}{3}$.

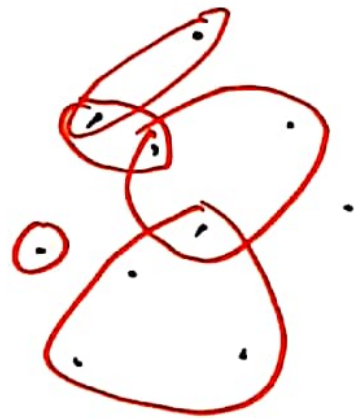
Let $A = \{s_i : rs_i \bmod p \in \mathcal{I}\}$. Then $|A| > \frac{n}{3}$

A is sum-free: if $s_i, s_j, s_k \in A$ with $s_i + s_j = s_k$

$$\text{then } rs_i + rs_j = rs_k$$

$$\text{X } \mathcal{I} \text{ sum-free mod } p \quad \square \quad \in \mathcal{I} \quad \in \mathcal{I} \quad \equiv \quad \in \mathcal{I} \quad \text{mod}(p)$$

A **hypergraph** H is an ordered pair (V, E)
where V is a (finite) set (of vertices) and E ,
the set of **hyperedges** is a set of ~~2-element~~
subsets of V .



Note E is a set:

eg. $\{1, 2, 3\}$ either
present or not.

Often $V = [n]$

Write 125 for

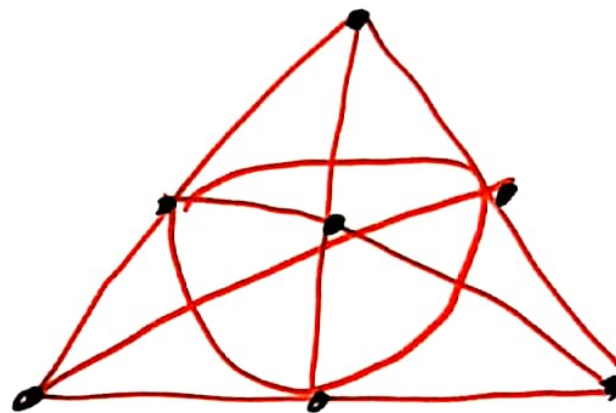
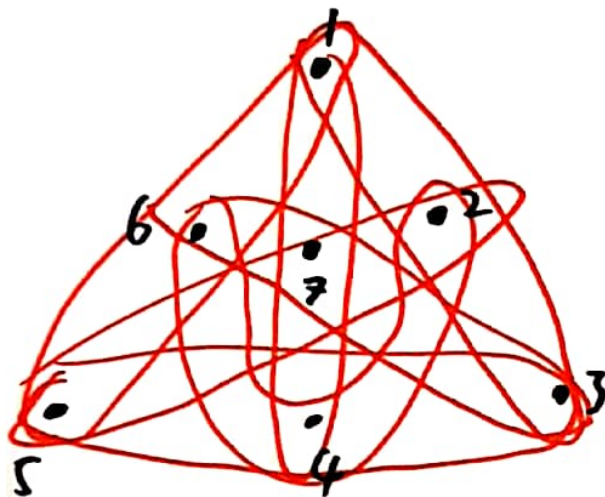
hyperedge $\{1, 2, 5\}$

(could consider multi-hypergraphs)
or

H is r -uniform if $|e| = r$ for all $e \in E(H)$

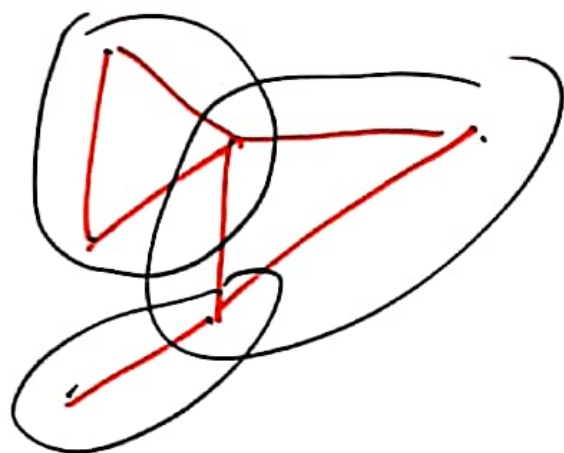
2-uniform hypergraph $\equiv$ graph.

eg. 3-uniform H : Fano plane

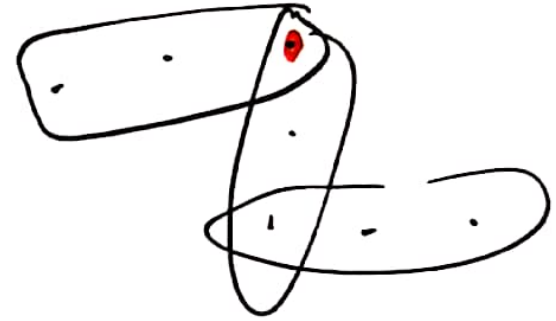
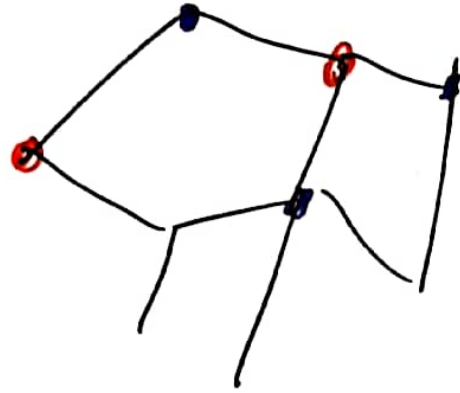


$V = [7]$ $E = \{123, 345, \dots$

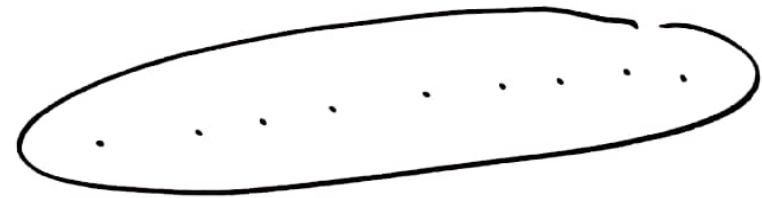
Colouring hypergraphs



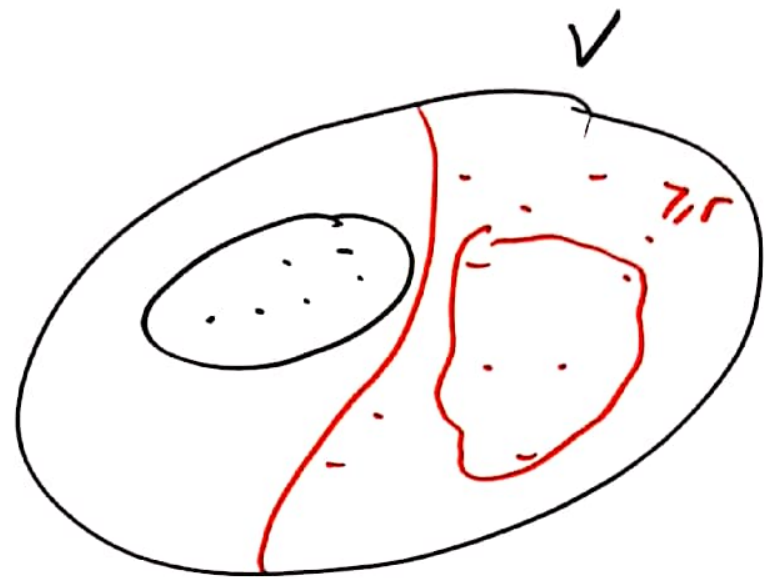
A proper k -colouring of a hypergraph H is a map $c: V(H) \rightarrow S$ where $|S| = k$ s.t. no edge e of H is monochromatic.



A hypergraph H is k -colourable if it has a proper k -colouring.



only forbid 2 combinations
out of 2^r .



complete r -uniform hypergraph on V

$E = \text{all } \binom{n}{r} \text{ } r\text{-element subsets } n = |V|$

If $n = 2r - 1$ not 2-colourable

$$e(H) = \binom{2r-1}{r} \approx 4^r.$$

Let $m(r) = \min \{e(H) : H \text{ is } r\text{-uniform \& NOT 2-colourable}\}$

$$m(2) = 3 \quad m(3) = 7 \quad (\text{Fano plane})$$

Theorem 1.6 For $r \geq 2$, $m(r) \geq 2^{r-1}$.

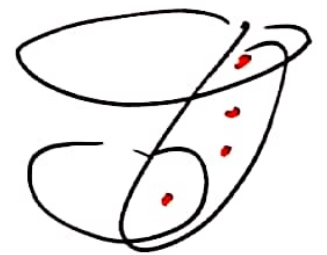


Pf. Let H be r -uniform with $e(H) = m < 2^{r-1}$.

Suffices to show H is 2-colourable.

Colouring vertices of H randomly, each red with prob. $\frac{1}{2}$, blue otherwise, independently.

Let A_i be event that i^{th} edge is mono.



$$IP(A_i) = 2/2^r. \quad \text{Union bound } IP(\cup A_i) \leq \sum IP(A_i) = \frac{m}{2^{r-1}} < 1$$

$IP(\text{colouring bad}) < 1. \therefore \exists \text{ good } \square \text{ colouring.}$

Theorem 1.7 (Erdős 1964). For r large enough $m(r) \leq 3r^2 2^r$.

Proof Fix $r \geq 3$

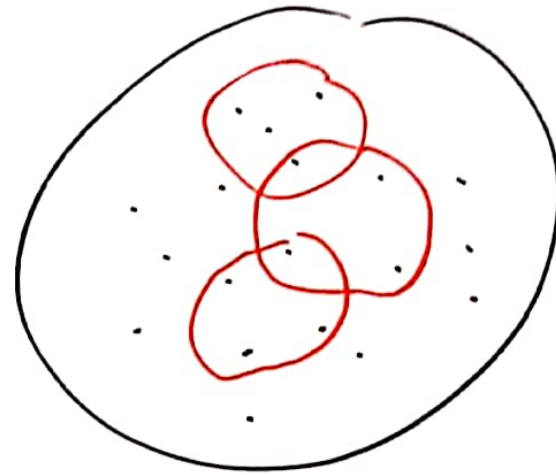
Let $m = 3r^2 2^r$

Task: show $\exists$ r -uniform H with $e(H) \leq m$ s.t. H
is not 2-colourable.

Idea: let H be random.

Fix a set of $n = r^2$ vertices

Pick $e_1, \dots, e_m$ independently
and uniformly at random from
all $\binom{n}{r}$ possible hyperedges.



$$|V| = n$$

Let $H = (V, \{e_1, \dots, e_m\})$

N.B. duplicate edges only 'there once'

$$\text{so } e(H) \leq m.$$

$$\{1, 1, 2\} = \{1, 2\}$$

Fix a (not random) 2-colouring of V .

Let $p = p_c = \mathbb{P}(e_i \text{ is c-mono})$

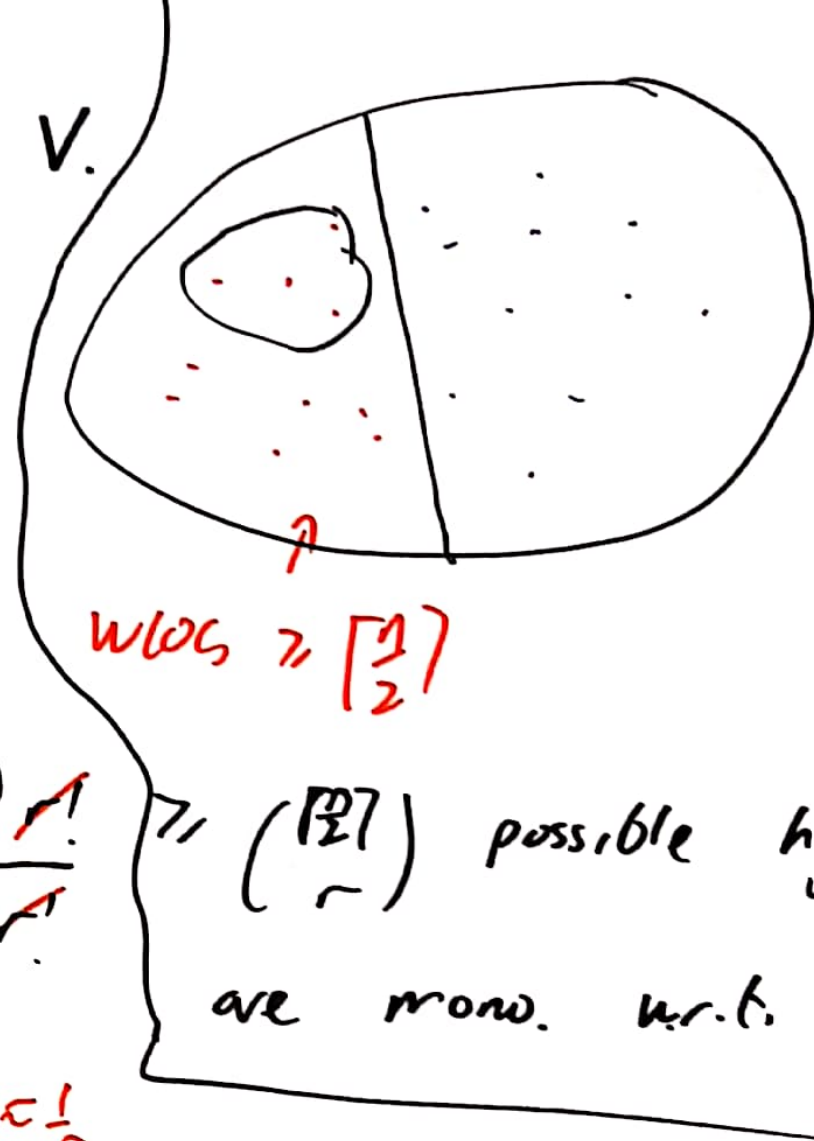
Then

$$p \geq \frac{\binom{\lceil \frac{n}{2} \rceil}{r}}{\binom{n}{r}} \geq \frac{(\frac{n}{2}) (\frac{n}{2} - 1) \dots (\frac{n}{2} - r + 1)}{n (n-1) \dots (n-r+1)}$$

ratios decrease

$$\geq \left(\frac{\frac{n}{2} - r}{n - r} \right)^r = \frac{1}{2^r} \left(1 - \frac{r}{n-r} \right)^r$$

$\frac{r}{n-r} \approx \frac{1}{r^2}$



$$\frac{n-1}{y-1} \quad \frac{n}{y}$$

$$P \geq \frac{1}{2^r} \left(1 - \frac{r}{r^2 - r}\right)^r = \frac{1}{2^r} \left(1 - \frac{1}{r-1}\right)^r \geq \left(\frac{1}{3 \cdot 2^r}\right)^r \stackrel{= p_0}{=} \text{if } r \geq r_0$$

assume
from now on

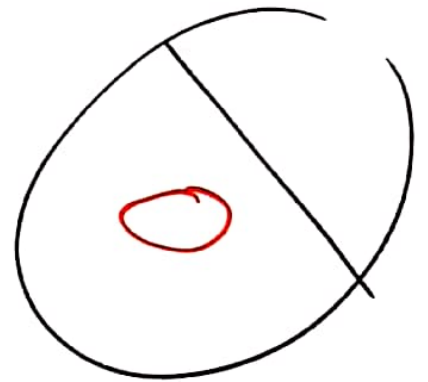
$$P(e_i \text{ is } c\text{-mono}) \geq p_0$$

$$P(c \text{ is a proper colouring of } H)$$

$$= P(\text{none of } e_1, \dots, e_m \text{ are } c\text{-mono})$$

$$= \prod_{i=1}^m P(e_i \text{ is not } c\text{-mono})$$

$$\leq (1 - p_0)^m$$



Let $A_c = \{c \text{ is a proper coloring of } H\}$

$$\mathbb{P}(H \text{ is 2-colourable}) = \mathbb{P}(\cup A_c) \leq \sum \mathbb{P}(A_c)$$

$$\leq 2^n (1-p_0)^m$$

$$\leq 2^n e^{-p_0 m}$$

$$1-x \leq e^{-x}$$

$$\leq 2^{r^2} e^{-\frac{m}{3 \cdot 2^r}} = 2^{r^2} e^{-r^2} < 1$$

H is always r -uniform & $e(H) \leq m$

Shown possible H is NOT 2-colourable.

□

$A_1, \dots, A_n$ events

IF disjoint $IP(\cup A_i) = \sum IP(A_i)$

IF indep. $IP(\cap A_i) = \prod IP(A_i)$ $IP(\cup A_i) = 1 - \prod_i (1 - IP(A_i))$

We ALWAYS have $IP(\cup A_i) \leq \sum IP(A_i)$

$$\{H \text{ is 2-colourable}\} = \bigcup_c \bigcap_i \{e_i \text{ is mono. wrt. } c\}^c$$