

Section V.3: Push-outs

Def G_0, G_1, G_2 groups, $\phi_1: G_0 \rightarrow G_1, \phi_2: G_0 \rightarrow G_2$ homom.
 $\langle X_1 | R_1 \rangle, \langle X_2 | R_2 \rangle$ canonical presentations of G_1, G_2 , resp.

The push-out $G_1 *_{G_0} G_2 := \langle X_1 \cup X_2 \mid R_1 \cup R_2 \cup \{\phi_1(g) = \phi_2(g) : g \in G_0\} \rangle$
 assume $X_1 \cap X_2 = \emptyset$

Rem $G_1 *_{G_0} G_2$ depends on ϕ_1, ϕ_2

Rem By Lemma V.11, $X_1 \hookrightarrow X_1 \cup X_2$ and $X_2 \hookrightarrow X_1 \cup X_2$
 induce homomorphisms $\kappa_1: G_1 \rightarrow G_1 *_{G_0} G_2, \kappa_2: G_2 \rightarrow G_1 *_{G_0} G_2$
 such that the following diagram commutes:

$$\begin{array}{ccc} G_0 & \xrightarrow{\phi_1} & G_1 \\ \phi_2 \downarrow & & \downarrow \kappa_1 \\ G_2 & \xrightarrow[\kappa_2]{} & G_1 *_{G_0} G_2 \end{array} \quad \kappa_1 \circ \phi_1 = \kappa_2 \circ \phi_2$$

Proposition V.19 (Universal property of push-outs)

$G_1 *_{G_0} G_2$ be the push-out of $G_0 \xrightarrow{\phi_1} G_1$
 $\phi_2 \downarrow$
 G_2

Let $\beta_1: G_1 \rightarrow H, \beta_2: G_2 \rightarrow H$ are group homom. such that

$G_0 \xrightarrow{\phi_1} G_1$ is commutative (i.e. $\beta_1 \circ \phi_1 = \beta_2 \circ \phi_2$).

$\phi_2 \downarrow$ $\downarrow \beta_1$

$G_2 \xrightarrow{\beta_2} H$

Then $\exists!$ homom $\beta: G_1 *_{G_0} G_2 \rightarrow H$ such that

$$\begin{array}{ccc} G_0 & \xrightarrow{\phi_1} & G_1 \\ \phi_2 \downarrow & & \downarrow \kappa_1 \\ G_2 & \xrightarrow[\kappa_2]{} & G_1 *_{G_0} G_2 \end{array} \quad \begin{array}{c} \beta_1 \\ \beta_2 \end{array} \quad \text{commutative (i.e. } \beta \circ \kappa_1 = \beta_1, \beta \circ \kappa_2 = \beta_2 \text{)}$$

Proof $G_1 *_{G_0} G_2$ is generated by $G_1 \cup G_2$.

Let $\beta(g_i) := \beta_i(g_i)$ for $g_i \in G_i, i \in \{1, 2\}$.

Forced on us by commutativity ($\beta \circ \kappa_1 = \beta_1, \beta \circ \kappa_2 = \beta_2$).

So, if β exists, it is unique.

β well-defined: $\beta(r) = e \quad \forall$ relation r for $G_1 *_{G_0} G_2$.

$\beta(r) = e$ if $r \in R_1 \cup R_2$, since β_1 and β_2 are well-defined
 The other type of relation is $a \phi_1(g) \phi_2^{-1}(a) = e, a \in G_0$.

But $\beta\phi_1(g)(\beta\phi_2(g))^{-1} = e$ by commutative square $(\beta_1 \circ \phi_1 = \beta_2 \circ \phi_2)$. \square

Lemma V.20. $\langle X'_1 | R'_1 \rangle, \langle X'_2 | R'_2 \rangle$ presentations of G_1, G_2 , resp.

$X'_1 \cap X'_2 = \emptyset$. Then $G_1 *_G G_2 \cong \langle X'_1 \cup X'_2 | R'_1 \cup R'_2 \cup \{\phi_1(g) = \phi_2(g) : g \in G_0\} \rangle$.

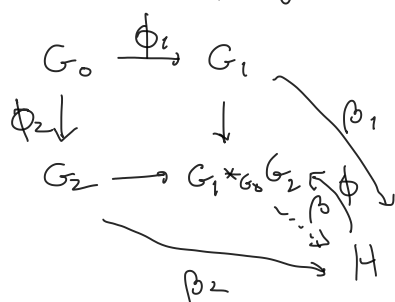
Proof $G := G_1 *_G G_2$

$G_1 \rightarrow \langle X'_1 | R'_1 \rangle, G_2 \rightarrow \langle X'_2 | R'_2 \rangle$ "identity" maps.

By lemma V.11, they induce homomorphisms

$\beta_1 : G_1 \rightarrow H$ and $\beta_2 : G_2 \rightarrow H$. Satisfy Prop. V.19 because

$\phi_1(g) = \phi_2(g) \forall g \in G_0$. So, by Prop. V.19, $\exists \beta : G \rightarrow H$ s.t. :



commutes. We define an inverse $\phi : H \rightarrow G$, as follows:

\exists function $X'_i \rightarrow G_i \xrightarrow{\alpha_i} G$

$f : X'_1 \cup X'_2 \rightarrow G$ induces a homom.

$\phi : F(X'_1 \cup X'_2) \rightarrow G$

By lemma V.11, this descends to a homom

$\phi : H \rightarrow G$ because $\phi(r) = e \in H \forall$ relation r in H .

ϕ is inverse of $\beta \Rightarrow G \cong H$. \square

Def $G_0 = \{e\}$, then $G_1 *_G G_2 = G_1 * G_2$ free product of G_1, G_2 (only depends on G_1, G_2).

Example $\mathbb{Z} * \mathbb{Z} \cong F\langle x, y \rangle$ because $\mathbb{Z} = \langle x | \rangle$

$\mathbb{Z} = \langle y | \rangle$, so $\mathbb{Z} * \mathbb{Z} = \langle x, y | \rangle$ by lemma V.20.

Def When $\phi_1 : G_0 \rightarrow G_1$ and $\phi_2 : G_0 \rightarrow G_2$ are injective, then we say that $G_1 *_G G_2$ amalgamated free product of G_1 and G_2 along G_0 .