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## Numerical Analysis Hilary Term 2024

### Lecture 11: Gauss quadrature

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**Terminology:** Quadrature  $\equiv$  numerical integration

**Goal:** given a (continuous) function  $f : [a, b] \rightarrow \mathbb{R}$ , find its integral  $I = \int_a^b f(x)dx$ , as accurately as possible.

**Idea: Approximate and Integrate.** Find a polynomial  $p_n$  from data  $\{(x_k, f(x_k))\}_{k=0}^n$  by Lagrange interpolation (lecture 1), and integrate  $\int_{x_0}^{x_n} p_n(x) dx =: I_n$ . Ideally,  $I_n = I$  or at least  $I_n \approx I$ . Is this true?

If we choose  $x_k$  to be equispaced points in  $[a, b]$ , the resulting  $I_n$  is known as the Newton-Cotes quadrature. This method is actually quite unstable and inaccurate, and a much more accurate and elegant quadrature rule exists: Gauss quadrature. In this lecture we cover this beautiful result involving orthogonal polynomials.

**Preparations:** Suppose that  $w$  is a weight function, defined, positive and integrable on the open interval  $(a, b)$  of  $\mathbb{R}$ .

**Lemma.** Let  $\{\phi_0, \phi_1, \dots, \phi_n, \dots\}$  be orthogonal polynomials for the inner product  $\langle f, g \rangle = \int_a^b w(x)f(x)g(x) dx$ . Then, for each  $k = 0, 1, \dots$ ,  $\phi_k$  has  $k$  distinct roots in the interval  $(a, b)$ .

**Proof.** Since  $\phi_0(x) \equiv \text{const.} \neq 0$ , the result is trivially true for  $k = 0$ . Suppose that  $k \geq 1$ :  $\langle \phi_k, \phi_0 \rangle = \int_a^b w(x)\phi_k(x)\phi_0(x) dx = 0$  with  $\phi_0$  constant implies that  $\int_a^b w(x)\phi_k(x) dx = 0$  with  $w(x) > 0$ ,  $x \in (a, b)$ . Thus  $\phi_k(x)$  must change sign in  $(a, b)$ , i.e.,  $\phi_k$  has at least one root in  $(a, b)$ .

Suppose that there are  $\ell$  points  $a < r_1 < r_2 < \dots < r_\ell < b$  where  $\phi_k$  changes sign for some  $1 \leq \ell \leq k$ . Then

$$q(x) = \prod_{j=1}^{\ell} (x - r_j) \times \text{the sign of } \phi_k \text{ on } (r_\ell, b)$$

has the same sign as  $\phi_k$  on  $(a, b)$ . Hence

$$\langle \phi_k, q \rangle = \int_a^b w(x)\phi_k(x)q(x) dx > 0,$$

and thus it follows from the previous lemma (cf. Lecture 12) that  $q$ , (which is of degree  $\ell$ ) must be of degree  $\geq k$ , i.e.,  $\ell \geq k$ . However,  $\phi_k$  is of exact degree  $k$ , and therefore the number of its distinct roots,  $\ell$ , must be  $\leq k$ . Hence  $\ell = k$ , and  $\phi_k$  has  $k$  distinct roots in  $(a, b)$ .  $\square$

**Application to quadrature.** The above lemma leads to very efficient quadrature rules since it answers the question: how should we choose the quadrature points  $x_0, x_1, \dots, x_n$  in the quadrature rule

$$\int_a^b w(x)f(x) dx \approx \sum_{j=0}^n w_j f(x_j) \tag{1}$$

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so that the rule is exact for polynomials of degree as high as possible? (The case  $w(x) \equiv 1$  is the most common.)

**Recall:** the Lagrange interpolating polynomial

$$p_n = \sum_{j=0}^n f(x_j) L_{n,j} \in \Pi_n$$

is unique, so  $f \in \Pi_n \implies p_n \equiv f$  whatever interpolation points are used, and moreover

$$\int_a^b w(x) f(x) dx = \int_a^b w(x) p_n(x) dx = \sum_{j=0}^n w_j f(x_j),$$

exactly, where

$$w_j = \int_a^b w(x) L_{n,j}(x) dx. \quad (2)$$

**Theorem.** Suppose that  $x_0 < x_1 < \dots < x_n$  are the roots of the  $n+1$ -st degree orthogonal polynomial  $\phi_{n+1}$  with respect to the inner product

$$\langle g, h \rangle = \int_a^b w(x) g(x) h(x) dx.$$

Then, the quadrature formula (1) with weights (2) is exact whenever  $f \in \Pi_{2n+1}$ .

**Proof.** Let  $p \in \Pi_{2n+1}$ . Then by the Division Algorithm  $p(x) = q(x)\phi_{n+1}(x) + r(x)$  with  $q, r \in \Pi_n$ . So

$$\int_a^b w(x) p(x) dx = \int_a^b w(x) q(x) \phi_{n+1}(x) dx + \int_a^b w(x) r(x) dx = \sum_{j=0}^n w_j r(x_j) \quad (3)$$

since the integral involving  $q \in \Pi_n$  is zero by the lemma above and the other is integrated exactly since  $r \in \Pi_n$ . Finally  $p(x_j) = q(x_j)\phi_{n+1}(x_j) + r(x_j) = r(x_j)$  for  $j = 0, 1, \dots, n$  as the  $x_j$  are the roots of  $\phi_{n+1}$ . So (3) gives

$$\int_a^b w(x) p(x) dx = \sum_{j=0}^n w_j p(x_j),$$

where  $w_j$  is given by (2) whenever  $p \in \Pi_{2n+1}$ . □

These quadrature rules are called **Gauss quadratures**.

- $w(x) \equiv 1$ ,  $(a, b) = (-1, 1)$ : Gauss–Legendre quadrature.
- $w(x) = (1 - x^2)^{-1/2}$  and  $(a, b) = (-1, 1)$ : Gauss–Chebyshev quadrature.
- $w(x) = e^{-x}$  and  $(a, b) = (0, \infty)$ : Gauss–Laguerre quadrature.
- $w(x) = e^{-x^2}$  and  $(a, b) = (-\infty, \infty)$ : Gauss–Hermite quadrature.

They give better accuracy than Newton–Cotes quadrature for the same number of function evaluations.

Note when using quadrature on unbounded intervals, the integral should be of the form  $\int_0^\infty e^{-x} f(x) dx$  and only  $f$  is sampled at the nodes.

Note that by the linear change of variable  $t = (2x - a - b)/(b - a)$ , which maps  $[a, b] \rightarrow [-1, 1]$ , we can evaluate for example

$$\int_a^b f(x) dx = \int_{-1}^1 f\left(\frac{(b-a)t + b + a}{2}\right) \frac{b-a}{2} dt \simeq \frac{b-a}{2} \sum_{j=0}^n w_j f\left(\frac{b-a}{2} t_j + \frac{b+a}{2}\right),$$

where  $\simeq$  denotes “quadrature” and the  $t_j$ ,  $j = 0, 1, \dots, n$ , are the roots of the  $n + 1$ -st degree Legendre polynomial.

**Example.** 2-point Gauss–Legendre quadrature:  $\phi_2(t) = t^2 - \frac{1}{3} \implies t_0 = -\frac{1}{\sqrt{3}}, t_1 = \frac{1}{\sqrt{3}}$  and

$$w_0 = \int_{-1}^1 \frac{t - \frac{1}{\sqrt{3}}}{-\frac{1}{\sqrt{3}} - \frac{1}{\sqrt{3}}} dt = - \int_{-1}^1 \left(\frac{\sqrt{3}}{2}t - \frac{1}{2}\right) dt = 1,$$

with  $w_1 = 1$ , similarly. So e.g., changing variables  $x = (t + 3)/2$ ,

$$\int_1^2 \frac{1}{x} dx = \frac{1}{2} \int_{-1}^1 \frac{2}{t+3} dt \simeq \frac{1}{3 + \frac{1}{\sqrt{3}}} + \frac{1}{3 - \frac{1}{\sqrt{3}}} = 0.6923077 \dots$$

Note that the trapezium rule (also two evaluations of the integrand) gives

$$\int_1^2 \frac{1}{x} dx \simeq \frac{1}{2} \left[ \frac{1}{2} + 1 \right] = 0.75,$$

whereas  $\int_1^2 \frac{1}{x} dx = \ln 2 = 0.6931472 \dots$

**Theorem.** Error in Gauss quadrature: suppose that  $f^{(2n+2)}$  is continuous on  $(a, b)$ . Then

$$\int_a^b w(x) f(x) dx = \sum_{j=0}^n w_j f(x_j) + \frac{f^{(2n+2)}(\eta)}{(2n+2)!} \int_a^b w(x) \prod_{j=0}^n (x - x_j)^2 dx,$$

for some  $\eta \in (a, b)$ .

**Proof.** The proof is based on the Hermite interpolating polynomial  $H_{2n+1}$  to  $f$  on  $x_0, x_1, \dots, x_n$ . [Recall that  $H_{2n+1}(x_j) = f(x_j)$  and  $H'_{2n+1}(x_j) = f'(x_j)$  for  $j = 0, 1, \dots, n$ .] The error in Hermite interpolation is

$$f(x) - H_{2n+1}(x) = \frac{1}{(2n+2)!} f^{(2n+2)}(\eta(x)) \prod_{j=0}^n (x - x_j)^2$$

for some  $\eta = \eta(x) \in (a, b)$ . Now  $H_{2n+1} \in \Pi_{2n+1}$ , so

$$\int_a^b w(x) H_{2n+1}(x) dx = \sum_{j=0}^n w_j H_{2n+1}(x_j) = \sum_{j=0}^n w_j f(x_j),$$

the first identity because Gauss quadrature is exact for polynomials of this degree and the second by interpolation. Thus

$$\begin{aligned} \int_a^b w(x)f(x) \, dx - \sum_{j=0}^n w_j f(x_j) &= \int_a^b w(x)[f(x) - H_{2n+1}(x)] \, dx \\ &= \frac{1}{(2n+2)!} \int_a^b f^{(2n+2)}(\eta(x))w(x) \prod_{j=0}^n (x - x_j)^2 \, dx, \end{aligned}$$

and hence the required result follows from the integral mean value theorem as  $w(x) \prod_{j=0}^n (x - x_j)^2 \geq 0$ .  $\square$

**Remark:** the “direct” approach of finding Gauss quadrature formulae sometimes works for small  $n$ , but more sophisticated algorithms are used for large  $n$ .<sup>1</sup>

**Example.** To find the two-point Gauss–Legendre rule  $w_0 f(x_0) + w_1 f(x_1)$  on  $(-1, 1)$  with weight function  $w(x) \equiv 1$ , we need to be able to integrate any cubic polynomial exactly, so

$$2 = \int_{-1}^1 1 \, dx = w_0 + w_1 \tag{4}$$

$$0 = \int_{-1}^1 x \, dx = w_0 x_0 + w_1 x_1 \tag{5}$$

$$\frac{2}{3} = \int_{-1}^1 x^2 \, dx = w_0 x_0^2 + w_1 x_1^2 \tag{6}$$

$$0 = \int_{-1}^1 x^3 \, dx = w_0 x_0^3 + w_1 x_1^3. \tag{7}$$

These are four nonlinear equations in four unknowns  $w_0, w_1, x_0$  and  $x_1$ . Equations (5) and (7) give

$$\begin{bmatrix} x_0 & x_1 \\ x_0^3 & x_1^3 \end{bmatrix} \begin{bmatrix} w_0 \\ w_1 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix},$$

which implies that

$$x_0 x_1^3 - x_1 x_0^3 = 0$$

for  $w_0, w_1 \neq 0$ , i.e.,

$$x_0 x_1 (x_1 - x_0)(x_1 + x_0) = 0.$$

If  $x_0 = 0$ , this implies  $w_1 = 0$  or  $x_1 = 0$  by (5), either of which contradicts (6). Thus  $x_0 \neq 0$ , and similarly  $x_1 \neq 0$ . If  $x_1 = x_0$ , (5) implies  $w_1 = -w_0$ , which contradicts (4). So  $x_1 = -x_0$ , and hence (5) implies  $w_1 = w_0$ . But then (4) implies that  $w_0 = w_1 = 1$  and (6) gives

$$x_0 = -\frac{1}{\sqrt{3}} \quad \text{and} \quad x_1 = \frac{1}{\sqrt{3}},$$

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<sup>1</sup>See e.g., the research paper by Hale and Townsend, “Fast and accurate computation of Gauss–Legendre and Gauss–Jacobi quadrature nodes and weights” SIAM J. Sci. Comput. 2013.

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which are the roots of the Legendre polynomial  $x^2 - \frac{1}{3}$ .

**Convergence:** Gauss quadrature converges astonishingly fast. It can be shown that if  $f$  is analytic on  $[a, b]$ , the convergence is geometric (exponential) in the number of samples. This is in contrast to other (more straightforward) quadrature rules:

- Newton-Cotes: Find interpolant in  $n$  equispaced points, and integrate interpolant. Convergence: (often) Divergent!
- (Composite) trapezium rule: Find piecewise-linear interpolant in  $n$  equispaced points, and integrate interpolant. Convergence:  $O(1/n^2)$  (assumes  $f''$  exists)
- (Composite) Simpson's rule: Find piecewise-quadratic interpolant in  $n$  equispaced points (each subinterval containing three points), and integrate interpolant. Convergence:  $O(1/n^4)$  (assumes  $f''''$  exists)

The figure below illustrates the performance on integrating the Runge function.

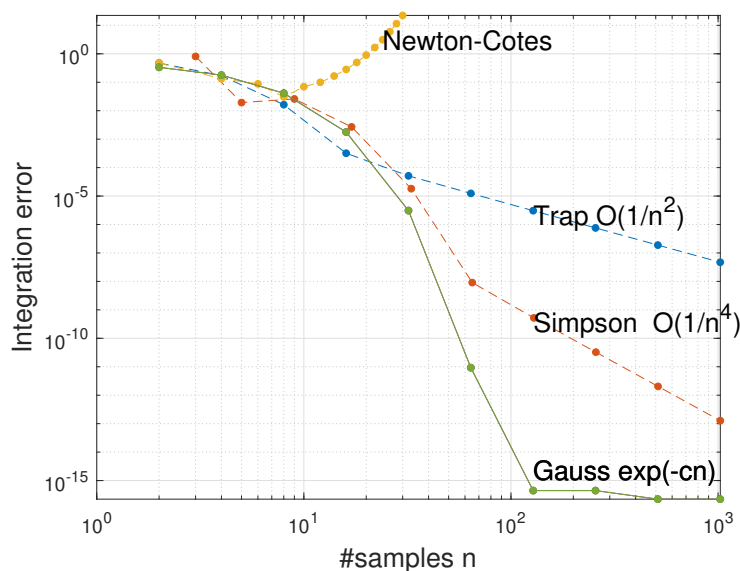
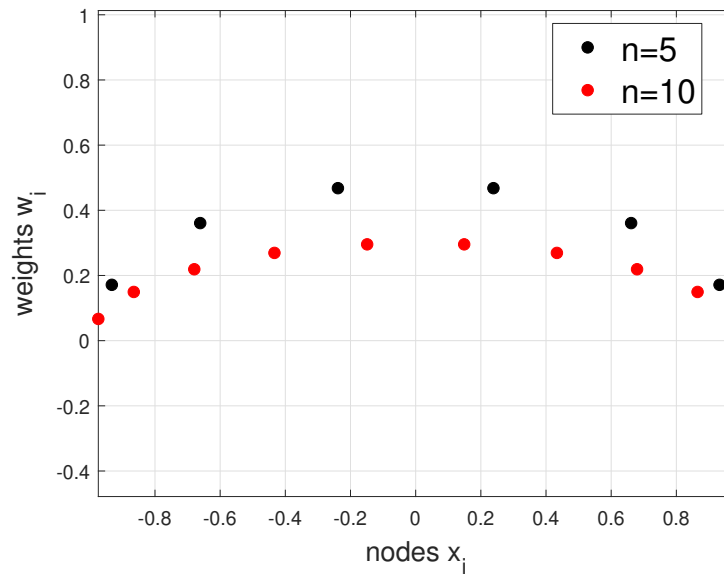


Figure 1: Convergence of quadrature rules for  $\int_{-1}^1 \frac{1}{25x^2+1} dx$  (Runge function)

**Nodes and weights for Gauss(-Legendre) quadrature** The figure below shows the nodes (interpolation points) and the corresponding weights with the standard Gauss-Legendre quadrature rule, i.e., when  $w(x) = 1$  and  $[a, b] = [-1, 1]$ . In Chebfun these are computed conveniently by `[x,w] = legpts(n+1)`



Note that the nodes/interpolation points cluster near endpoints (and sparser in the middle); this is a general phenomenon, and very analogous to the Chebyshev interpolation points mentioned in the least-squares lecture (Gauss and Chebyshev points have asymptotically the same distribution of points). Note also that the weights are all positive and shrink as  $n$  grows; they have to because they sum to 2 (why?).