

## Geometric Group Theory

### Problem Sheet 4

We use the notation from Lecture Notes,  $X \sim Y$ , for two metric spaces that are quasi-isometric.

1. i) Show that the relation of quasi-isometry of metric spaces  $\sim$  is an equivalence relation.

ii) Let  $S_1, S_2$  be finite generating sets of a group  $G$ . Show that  $\Gamma(S_1, G) \sim \Gamma(S_2, G)$ .

*Solution.* i) Let  $f : X \rightarrow Y$  a  $(K, A)$ -quasi-isometry. Define a 'quasi-inverse'  $g : Y \rightarrow X$  as follows: Given  $y \in Y$  pick  $x \in X$  such that  $d(y, f(x)) \leq A$ . Define  $g(y) = x$ . Then  $g$  is also a quasi-isometry: Let  $x \in X$  and  $y = f(x)$  then  $g(y) = x_1$  for some  $x_1$  for which  $d(f(x), f(x_1)) \leq A$ . So  $d(x, x_1) \leq KA + A$ .

ii) We consider the identity map on the vertices  $f : \Gamma(G, S_1) \rightarrow \Gamma(G, S_2)$ . We can write each element of  $S_1$  as a word on  $S_2$  and each element of  $S_2$  as a word on  $S_1$ . The maximum length of all these words controls the quasi-isometry constants.

2. Given  $\epsilon, \delta > 0$  a subset  $N$  of a metric space  $X$  is called an  $(\epsilon, \delta)$ -net (or simply a net) if for every  $x \in X$  there is some  $n \in N$  such that  $d(x, n) \leq \epsilon$  and for every  $n_1, n_2 \in N$ ,  $d(n_1, n_2) \geq \delta$ .

A set  $N$  that satisfies only the second condition (i.e. for every  $n_1, n_2 \in N$ ,  $d(n_1, n_2) \geq \delta$ ) is called  $\delta$ -separated.

i) Show that any metric space  $X$  has a  $(1, 1)$ -net.

ii) Show that if  $N \subset X$  is a net then  $X \sim N$ .

iii) Show that  $X \sim Y$  if and only if there are nets  $N_1 \subset X, N_2 \subset Y$  and a bilipschitz map  $f : N_1 \rightarrow N_2$ .

iv) Let  $G$  be a f.g. group. Show that  $H < G$  is a net in  $G$  if and only if  $H$  is a finite index subgroup of  $G$ .

*Solution.* i) Let  $N$  be a maximal subset of  $X$  such that for any  $a, b \in N$   $d(a, b) \geq 1$ . Such an  $N$  exists by Zorn's lemma. Now if  $x \in X$  and  $d(x, a) \geq 1$  for any  $a \in N$  then  $N$  is not maximal. So there is some  $a \in N$  such that  $d(a, x) \leq 1$ .

ii) The inclusion  $N \rightarrow X$  is a quasi-isometry.

iii) Let  $f : X \rightarrow Y$  be a  $(K, A)$ -quasi-isometry. Pick  $N_1$  an  $(n, n)$ -net in  $X$  with  $n = 2K(A + 1) + A$  (sufficiently large). Then  $d(f(x), f(y)) \geq 1$  for  $x \neq y$  so  $f$  is injective on  $N_1$ . Also

$$d(f(x), f(y)) \leq Kd(x, y) + A \leq KAd(x, y)$$

and

$$d(f(x), f(y)) \geq \frac{d(x, y)}{K} - A \geq \frac{d(x, y)}{K} - \frac{d(x, y)}{2K} \geq \frac{d(x, y)}{2K}$$

Finally since for any  $y \in Y$  there is an  $x \in X$  such that  $d(y, f(x)) \leq A$  and there is an  $a \in N_1$  with  $d(a, x) \leq n$  we have that

$$d(f(n), y) \leq A + Kn + A.$$

so  $f(N_1) = N_2$  is a net in  $Y$ .

iv) Clearly if  $H$  is of index  $n$  then  $H$  is an  $(n, 1)$  net in  $G$ . Assume that  $H$  is an  $(n, 1)$  net in  $G$ . Let's say that there are  $M$  words on the generating set of  $G$  of length  $\leq n$ . For every  $g \in G$   $gw \in H$  for some word of length  $\leq n$ . So  $g \in Hw^{-1}$ . It follows that the index of  $H$  in  $G$  is bounded by  $M$ .

**3.** Prove that for every  $K \geq 1$  and  $A \geq 0$  there exists  $\lambda \geq 1$ ,  $\mu \geq 0$  and  $D \geq 0$  such that the following is true. Given a  $(K, A)$ -quasi-geodesic  $q : I \rightarrow X$  of endpoints  $x, y$  in a geodesic metric space  $X$  there exists a (continuous) path  $\alpha : I' \rightarrow X$  of endpoints  $x, y$  such that:

1. for all  $t, s \in I$ ,

$$\text{length}(\alpha([t, s])) \leq \lambda d(\alpha(t), \alpha(s)) + \mu;$$

2. for every  $x \in I$ ,  $d(q(x), \alpha(I')) \leq D$ ;

3. for every  $t \in I'$ ,  $d(\alpha(t), q(I)) \leq D$ .

*Solution.* Let  $t_0, t_1, \dots, t_n$  be points in the interval  $I$  such that  $t_0, t_n$  are its endpoints,  $|t_{i+1} - t_i| = 1$  for all  $0 \leq i \leq n-1$ ,  $|t_{n+1} - t_n| \leq 1$ . Consider  $\alpha$  to be the polygonal line with geodesic edges  $[x, q(t_1)] \cup [q(t_1), q(t_2)] \cup \dots \cup [q(t_{n-1}), y]$ , parametrized by its arc length.

The last two conditions are satisfied with  $D = \frac{K}{2} + A$  and the first with  $\lambda = K^2$  and  $\mu = K(A + 1) + 2A$ .

**4.** Let  $X$  be a  $\delta$ -hyperbolic geodesic metric space. If  $L$  is a geodesic in  $X$  and  $a \in X$  we say that  $b \in L$  is a projection of  $a$  to  $L$  if

$$d(a, b) = \inf\{d(a, x) : x \in L\}.$$

Show that if  $b_1, b_2$  are projections of  $a$  to  $L$  then  $d(b_1, b_2) \leq 2\delta$ .

*Solution.* This follows easily by considering the geodesic triangle  $[a, b_1, b_2]$ .

**5.** Let  $X$  be a geodesic metric space.

If  $\Delta = [x, y, z]$  is a geodesic triangle in  $X$ , then there is a metric tree (a 'tripod' if  $\Delta$  is not degenerate)  $T_\Delta$  with vertices  $x', y', z'$  (the endpoints when  $T_\Delta$  is not a segment) such that there is an onto map  $f_\Delta : \Delta \rightarrow T_\Delta$  that

restricts to an isometry from each side  $[x, y], [y, z], [x, z]$  to the corresponding segments  $[x', y'], [y', z'], [x', z']$  in the tree. We denote by  $c_\Delta$  the point  $[x', y'] \cap [y', z'] \cap [x', z']$  of  $T_\Delta$ .

We say that a geodesic triangle  $\Delta = [x, y, z]$  in a geodesic metric space is  $\delta$ -thin if for every  $t \in T_\Delta = [x', y', z']$ ,  $\text{diam}(f_\Delta^{-1}(t)) \leq \delta$ .

Prove that the following are equivalent:

1. There is a  $\delta \geq 0$  such that all geodesic triangles in  $X$  are  $\delta$ -slim.
2. There is a  $\delta' \geq 0$  such that all geodesic triangles in  $X$  are  $\delta'$ -thin.

*Solution.* This appears as Theorem 6.4, with proof, in the Lecture Notes. Please make sure that in class the students understand the two definitions and their equivalence.

**6.** Let  $G = \langle S \rangle$  be  $\delta$ -hyperbolic for some  $\delta \in \mathbb{N}$ ,  $\delta \geq 1$ .

1. Assume that for some  $g \in G, x \in \Gamma(S, G)$  with  $d(x, gx) > 100\delta$  we have that  $d(x, g^2x) \geq 2d(x, gx) - 12\delta$ .

Prove that

$$d(x, g^n x) \geq nd(x, gx) - 16n\delta$$

for all  $n \in \mathbb{N}$ .

2. Assume that  $g$  is an element of infinite order in  $G$ . Prove that there are constants  $c > 0, d \geq 0$  such that

$$d(1, g^n) \geq cn - d$$

for all  $n \in \mathbb{N}$ .

3. Show that  $G$  has no subgroup isomorphic to  $\langle x, t | txt^{-1} = x^2 \rangle$ .

*Solution.* 1. This is Lemma 6.4 in the Lecture Notes.

2. This is Proposition 6.4 in the Lecture Notes.

3.  $t^n xt^{-n} = x^{2^n}$  which contradicts the fact that  $x^n$  is a quasi-geodesic.

**7.** Let  $G = \langle S | R \rangle$  be a Dehn presentation of a  $\delta$ -hyperbolic group. Show that we can decide whether a word  $w$  on  $S$  represents an infinite order element.

*Solution.* To clarify, our input for the algorithm is the finite presentation  $\langle S | R \rangle$  and  $\delta$ .

*1st solution:* We use a Dehn presentation and using the solution to the conjugacy problem we check successively for the powers of  $w$ ,  $w^k$ , whether they are conjugate to an element of length  $\leq \max\{|r| + 2\}$  where  $r$  ranges

over all relations of the Dehn presentation. Eventually we will either find that  $w^k = 1$  or we will find two powers  $w^k, w^m$  which are conjugate to the same element  $a$ . It follows that these are conjugate so there is some  $t$  such that  $tw^kt^{-1} = w^m$ . However this contradicts the fact that  $\langle w \rangle$  is a quasi-geodesic as in exercise 8. So either some power is equal to 1 or some power is not conjugate to any element of length  $\leq \max\{|r| + 2\}$  (and hence  $w$  is of infinite order).

*2nd solution:* Enumerate powers  $w^n$  and check if they are equal to 1. In parallel try to find a vertex  $m$  of the Cayley graph and a power  $w^k$  such that  $d(w^{2k}m, w^km) > 2d(m, w^km) - 12\delta$  and  $d(e, w^k) > 100\delta$ . If  $w$  is of finite order the first procedure will terminate. If  $w$  is of infinite order then by the proof of the proposition 6.4 in the notes showing that  $\langle w \rangle$  is a quasi-geodesic  $w^k$  and  $m$  with the above properties exist and we can detect them since the word problem is solvable in  $G$ .

**8.** Let  $G = \langle S | R \rangle$  be a Dehn presentation of a  $\delta$ -hyperbolic group. Show that we can decide whether a word  $w$  on  $S$  lies in the subgroup  $\langle v \rangle$ .

*Solution.* To clarify, our input for the algorithm is the finite presentation  $\langle S | R \rangle$ ,  $\delta$  and the words  $v, w$ .

The proof of proposition 6.4 shows that there is some vertex  $m$  in the Cayley graph and some power  $v^k$  such that  $d(v^{2k}m, v^km) \geq 2d(v^km, m) - 12\delta$ . However since we can solve the word problem we can find  $v^k, m$  just by calculating multiplication tables for larger and larger balls and powers of  $v$ . Once those are found we get an estimate, as in proposition 6.4, of the form  $d(v^n, e) \geq cn - d$  for some  $c, d > 0$ . So it is enough to check whether  $c^n = w$  for all  $n$  for which  $cn - d \leq |w|$ .