

The exponential function

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In *Analysis II* (previewed in *Analysis I*) one encounters the following proof of

$$\exp(x + y) = \exp(x) \exp(y), \tag{1}$$

where $\exp(x)$ is defined by

$$\exp(x) := \sum_{n=0}^{\infty} \frac{x^n}{n!}. \tag{2}$$

Proof. First we show that (2) has infinite radius of convergence (by, e.g., the *Ratio Test*). Then apply the *Differentiation Theorem for power series* to differentiate term by term:

$$\frac{d}{dx} \exp(x) = \sum_{n=0}^{\infty} \frac{d}{dx} \frac{x^n}{n!} = \sum_{n=1}^{\infty} \frac{x^{n-1}}{(n-1)!} = \sum_{n=0}^{\infty} \frac{x^n}{n!} = \exp(x).$$

Now differentiate $F(x) := \exp(x) \exp(c - x)$ with respect to x using the *Product Rule* and *Chain Rule* to get

$$\begin{aligned} \frac{d}{dx} F(x) &= \exp'(x) \exp(c - x) + \exp(x) \exp'(c - x)(-1) \\ &= \exp(x) \exp(c - x) - \exp(x) \exp(c - x) = 0. \end{aligned}$$

Now use the *Constancy Theorem* to deduce that the function $F(x)$ is constant on \mathbb{R} , so $\exp(x) \exp(c - x) = F(x) = F(0) = \exp(0) \exp(c)$. Substituting $c = x + y$ and noting that $\exp(0) = 1$ now gives the result (1). \square

There are a couple of things that are rather unsatisfactory about this proof.

1. It uses a lot of machinery from *Analysis II* to prove something that is very simple, and certainly should be accessible using just the methods of *Analysis I*.
2. It only works for *real* x and y . To generalise to complex numbers needs significant extra effort.

To be fair, the conclusion of the *Constancy Theorem* also applies to complex functions, and indeed much more will be proved in the Part A course *Metric Spaces and Complex Analysis*. But in some sense that makes the 1st point above even worse — one needs to wait until the 2nd year to see a complete proof of this basic identity.

Here we present two proofs of (1) that rely only on results from *Analysis I* and also work for complex numbers. The first is in some sense more natural and elementary¹, but uses a different definition of the exponential. The second uses (2), but is slightly more opaque. We then show the two definitions give the same function.

For the first proof we try to use the following definition².

$$\exp(x) = \lim_{n \rightarrow \infty} \left(1 + \frac{x}{n}\right)^n.$$

The hard part is showing that this converges. For *real* x it is not too difficult. A slight generalisation of question 4 on problem sheet 1 of *Analysis I* shows that the sequence is increasing in n for $n > |x|$ and bounded above. It is a bit more tricky for complex x however. To make life a little easier, we will just show convergence along a certain subsequence. This will be enough and makes the algebra a bit simpler. Hence we will define (using a different notation for \exp as we haven't shown it matches (2) yet)

$$e(x) := \lim_{n \rightarrow \infty} e_n(x), \quad \text{where} \quad e_n(x) := \left(1 + \frac{x}{2^n}\right)^{2^n}. \quad (3)$$

We first observe that $e_n(x)$ satisfies the following easy identity:

$$e_{n+1}(x) = \left(1 + \frac{x/2}{2^n}\right)^{2^n \cdot 2} = e_n\left(\frac{x}{2}\right)^2. \quad (4)$$

The following lemma shows that $e_n(x) \approx 1 + x$ for x small with an explicit error bound.

Lemma 1. For $|x| \leq \frac{1}{2}$ and any $n \geq 0$, $|e_n(x) - (1 + x)| \leq |x|^2$.

Proof. We use induction on n , $n = 0$ is trivial as $e_0(x) = 1 + x$. Now for $|x| \leq \frac{1}{2}$ and $n \geq 0$ write $e_n(\frac{x}{2}) = 1 + \frac{x}{2} + \eta$ where, by induction, $|\eta| \leq |\frac{x}{2}|^2 = \frac{|x|^2}{4}$. Then

$$\begin{aligned} |e_{n+1}(x) - (1 + x)| &= |e_n\left(\frac{x}{2}\right)^2 - (1 + x)| && \text{from (4)} \\ &= \left|(1 + \frac{x}{2} + \eta)^2 - (1 + x)\right| && e_n\left(\frac{x}{2}\right) = 1 + \frac{x}{2} + \eta \\ &= \left|\frac{x^2}{4} + \eta(2 + x + \eta)\right| && \text{expand and simplify} \\ &\leq \frac{|x|^2}{4} + |\eta|(2 + |x| + |\eta|) && \text{triangle inequality} \\ &\leq \frac{|x|^2}{4} + \frac{|x|^2}{4}(2 + |x| + \frac{|x|^2}{4}) && |\eta| \leq \frac{|x|^2}{4} \\ &\leq \frac{|x|^2}{4}\left(1 + 2 + \frac{1}{2} + \frac{1}{16}\right) && |x| \leq \frac{1}{2} \\ &\leq |x|^2. && \square \end{aligned}$$

The next step is the crucial bit: showing $e_n(x)$ converges.

¹'Elementary' is a technical term here. It means the proof does not use advanced concepts and theorems, but instead uses more basic techniques. It does *not* however mean the proof is simple or easy!

²Think 'continuous compound interest'.

Lemma 2. For any fixed $x \in \mathbb{C}$, $e_n(x)$ converges as $n \rightarrow \infty$.

Proof. We first note that for any $n \geq 0$ and $|x| \leq \frac{1}{2}$, $|e_n(x)| \leq 1 + |x| + |x|^2 \leq \frac{7}{4}$ by Lemma 1 and the triangle inequality. Now assume $|x| \leq \frac{1}{2}$. We show that for all $n > m \geq 0$,

$$|e_n(x) - e_m(x)| \leq \left(\frac{7}{8}\right)^m |x|^2 \quad (5)$$

by induction on m (simultaneously for all $n > m$). Lemma 1 is just the case $m = 0$ and, assuming the result for m ,

$$\begin{aligned} |e_{n+1}(x) - e_{m+1}(x)| &= \left| e_n\left(\frac{x}{2}\right)^2 - e_m\left(\frac{x}{2}\right)^2 \right| \\ &= \left| e_n\left(\frac{x}{2}\right) + e_m\left(\frac{x}{2}\right) \right| \left| e_n\left(\frac{x}{2}\right) - e_m\left(\frac{x}{2}\right) \right| \\ &\leq \left(\frac{7}{4} + \frac{7}{4}\right) \cdot \left(\frac{7}{8}\right)^m \left|\frac{x}{2}\right|^2 = \left(\frac{7}{8}\right)^{m+1} |x|^2 \end{aligned}$$

for all $n + 1 > m + 1$. Hence (5) holds for all $n > m \geq 0$.

Now $\left(\frac{7}{8}\right)^m |x|^2 \rightarrow 0$ as $m \rightarrow \infty$, so $(e_n(x))$ is a Cauchy sequence. Thus by the *Cauchy Convergence Criterion*, $(e_n(x))$ converges as $n \rightarrow \infty$ for any $x \in \mathbb{C}$ with $|x| \leq \frac{1}{2}$.

For larger $|x|$ we note that if $e_n(x)$ converges as $n \rightarrow \infty$ for any $|x| \leq K$ then $e_n(x) = e_{n-1}\left(\frac{x}{2}\right)^2$ converges for any $|x| \leq 2K$ by AOL. Hence we can inductively show $(e_n(x))$ converges as $n \rightarrow \infty$ for $|x| \leq \frac{2^t}{2}$ for all $t \geq 0$. But for any $x \in \mathbb{C}$, $|x| \leq \frac{2^t}{2}$ for some t . Thus $(e_n(x))$ converges as $n \rightarrow \infty$ for all $x \in \mathbb{C}$. \square

Having shown $e(x)$ is well defined, the main result is relatively straightforward.

Theorem 3. For all $x, y \in \mathbb{C}$, $e(x + y) = e(x)e(y)$.

Proof. We observe that

$$\frac{e_n(x)e_n(y)}{e_n(x+y)} = \left(\frac{(1 + \frac{x}{2^n})(1 + \frac{y}{2^n})}{1 + \frac{x+y}{2^n}} \right)^{2^n} = \left(1 + \frac{\frac{xy}{4^n}}{1 + \frac{x+y}{2^n}} \right)^{2^n} = \left(1 + \frac{z_n}{2^n} \right)^{2^n} = e_n(z_n), \quad (6)$$

where $z_n = \frac{xy/2^n}{1+(x+y)/2^n}$ (and the denominator in (6) is nonzero for sufficiently large n). But $z_n \rightarrow 0$ as $n \rightarrow \infty$ by the fact that $2^{-n} \rightarrow 0$ and AOL so, for sufficiently large n , $|z_n| \leq \frac{1}{2}$. But then $|e_n(z_n) - (1 + z_n)| \leq |z_n|^2$ by Lemma 1. But $z_n \rightarrow 0$, so $e_n(z_n) \rightarrow 1$ by sandwiching.

Now take (6) in the form $e_n(x+y)e_n(z_n) = e_n(x)e_n(y)$ and apply AOL and Lemma 2 to both sides to obtain $e(x+y) = e(x)e(y)$. \square

We remark that this proof uses the *Cauchy Convergence Criterion*, which was pretty much inevitable as we wanted to show a sequence of complex numbers converged to something we could not previously describe. (In the other proofs it is hidden away in the various results that are used.) Other than that it just uses AOL, sandwiching, and simple bounding techniques (triangle inequality and standard algebra).

We also remark that once differentiation is introduced, Lemma 1 tells us precisely that $e'(0) = 1$. Then $e'(x) = e(x)$ can be deduced from the *Chain Rule* and Theorem 3: $e(c)e'(x) = \frac{d}{dx}e(c)e(x) = \frac{d}{dx}e(x+c) = e'(x+c) \cdot 1$, and set $x = 0$.

We now come to the second proof of (1). This uses more properties of series from *Analysis I*. We start with a general result that could have been proved in *Analysis I* and is of independent interest (and was question 3 on the 2013 Prelims MII exam!).

Lemma 4 (Multiplication of absolutely convergent series). *If $\sum a_k$ and $\sum b_k$ converge absolutely and we define $c_k := a_0b_k + a_1b_{k-1} + \dots + a_kb_0 = \sum_{i=0}^k a_ib_{k-i}$, then $\sum c_k$ converges absolutely and*

$$\sum_{k=0}^{\infty} c_k = \sum_{k=0}^{\infty} a_k \sum_{k=0}^{\infty} b_k.$$

We remark that this is basically a theorem about rearranging terms in a series, summing them up in a different order. As a result it should not come as much of a surprise that the condition of *absolute* convergence is needed. The lemma would be false in general if the sequences converged, but not absolutely.

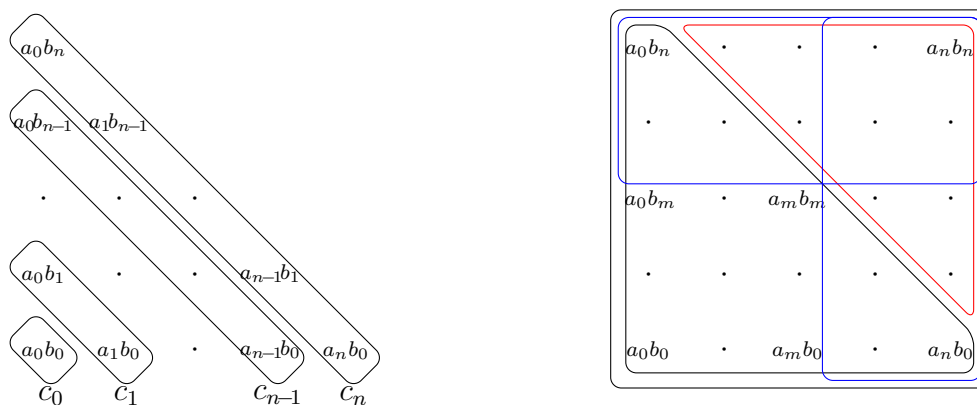
Before writing out a formal proof of this lemma, let's give the general idea of the proof. We want to relate

$$\sum_{k=0}^n c_k = \sum_{k=0}^n \sum_{i=0}^k a_ib_{k-i}$$

to

$$\sum_{i=0}^n a_i \sum_{j=0}^n b_j = \sum_{i=0}^n \sum_{j=0}^n a_ib_j$$

for some large n and let $n \rightarrow \infty$. If we consider both sums as sums of terms of the form a_ib_j and arrange these terms in a grid, the first sum is summing over a triangle of terms, while the second is summing over a square.



We bound the difference (the sum of terms in the red triangle) by showing that the sum of the absolute values outside a large square is small, noting that the original triangle does contain a square of size $m \approx n/2$. More specifically, we show that the sum in each blue rectangle is small.

Proof. By absolute convergence of $\sum a_k$ and $\sum b_k$ we can define $A := \sum_{k=0}^{\infty} |a_k|$ and $B := \sum_{k=0}^{\infty} |b_k|$. Now fix $\varepsilon > 0$. By convergence of the sums defining A and B , there exists an m such that $\sum_{k=m+1}^{\infty} |a_k| < \frac{\varepsilon}{A+B}$ and $\sum_{k=m+1}^{\infty} |b_k| < \frac{\varepsilon}{A+B}$. Now

$$\begin{aligned} \sum_{i=0}^n a_i \sum_{j=0}^n b_j - \sum_{k=0}^n c_k &= \sum_{i=0}^n \sum_{j=0}^n a_i b_j - \sum_{k=0}^n \sum_{i+j=k} a_i b_j \\ &= \sum_{i=0}^n \sum_{j=0}^n a_i b_j - \sum_{\substack{i=0 \\ i+j \leq n}}^n \sum_{j=0}^n a_i b_j = \sum_{i=0}^n \sum_{\substack{j=0 \\ i+j > n}}^n a_i b_j. \end{aligned}$$

We fix $n > 2m$. Then if $i + j > n$ either $i > m$ or $j > m$. Thus by the triangle inequality

$$\begin{aligned} \left| \sum_{i=0}^n a_i \sum_{j=0}^n b_j - \sum_{k=0}^n c_k \right| &\leq \sum_{\substack{i=0 \\ i+j > n}}^n \sum_{j=0}^n |a_i b_j| \leq \sum_{i=m+1}^n \sum_{j=0}^n |a_i b_j| + \sum_{i=0}^n \sum_{j=m+1}^n |a_i b_j| \\ &= \sum_{i=m+1}^n |a_i| \sum_{j=0}^n |b_j| + \sum_{i=0}^n |a_i| \sum_{j=m+1}^n |b_j| \\ &< \frac{\varepsilon}{A+B} \cdot B + A \cdot \frac{\varepsilon}{A+B} = \varepsilon. \end{aligned}$$

As for any $\varepsilon > 0$ we can find an m such that this holds for all $n > 2m$, we have that $|\sum_{k=0}^n a_k \sum_{k=0}^n b_k - \sum_{k=0}^n c_k| \rightarrow 0$ as $n \rightarrow \infty$. By AOL and sandwiching, $\sum c_k$ converges and $\sum_{k=0}^{\infty} c_k = \sum_{k=0}^{\infty} a_k \sum_{k=0}^{\infty} b_k$.

For absolute convergence, replace a_k and b_k by $a'_k := |a_k|$ and $b'_k := |b_k|$ and note that $\sum c'_k$ converges where $c'_k := \sum |a_i b_{k-i}| \geq |c_k|$. Now $\sum |c_k|$ converges by the *Comparison Test*. \square

Theorem 5. For all $x, y \in \mathbb{C}$, $\exp(x + y) = \exp(x) \exp(y)$.

Proof. For any $x, y \in \mathbb{C}$ the power series for $\exp(x)$ and $\exp(y)$ converge absolutely as the series has infinite radius of convergence (by e.g., the *Ratio Test*). Hence by Lemma 4

$$\begin{aligned} \exp(x) \exp(y) &= \sum_{n=0}^{\infty} \frac{x^n}{n!} \sum_{n=0}^{\infty} \frac{y^n}{n!} = \sum_{n=0}^{\infty} \sum_{i=0}^n \frac{x^i}{i!} \cdot \frac{y^{n-i}}{(n-i)!} \\ &= \sum_{n=0}^{\infty} \frac{1}{n!} \sum_{i=0}^n \binom{n}{i} x^i y^{n-i} = \sum_{n=0}^{\infty} \frac{(x+y)^n}{n!} = \exp(x+y), \end{aligned}$$

where in the second line we have used the definition $\binom{n}{i} = \frac{n!}{i!(n-i)!}$ and the *Binomial Theorem*. \square

Now we show that the $e(x)$ defined earlier is actually the same as $\exp(x)$.

Theorem 6. For all $x \in \mathbb{C}$, $(1 + \frac{x}{n})^n \rightarrow \exp(x)$ as $n \rightarrow \infty$. In particular $e(x) = \exp(x)$.

Direct proof without assuming Theorems 3 and 5. Use the *Binomial Theorem*

$$\left(1 + \frac{x}{n}\right)^n = 1 + \binom{n}{1} \left(\frac{x}{n}\right) + \binom{n}{2} \left(\frac{x}{n}\right)^2 + \cdots + \binom{n}{n} \left(\frac{x}{n}\right)^n. \quad (7)$$

We note that for each *fixed* k

$$\binom{n}{k} \left(\frac{x}{n}\right)^k = \frac{n(n-1)\cdots(n-k+1)}{n^k \cdot k!} x^k = \frac{x^k}{k!} \prod_{j=0}^{k-1} \left(1 - \frac{j}{n}\right) \rightarrow \frac{x^k}{k!}, \quad (8)$$

as $n \rightarrow \infty$ by AOL. But we are not done yet! We can't apply AOL to the whole expression (7) as the number of terms grows with n . Instead we show that terms for large k are small enough that we can throw away most of the terms and only lose at most some small ε .

To do this we note that in (8) we actually always have $|\binom{n}{k} (\frac{x}{n})^k| \leq |\frac{x^k}{k!}|$, so each term in the binomial expansion of $(1 + \frac{x}{n})^n$ is no larger than the corresponding term in $\exp(x)$.

So fix $\varepsilon > 0$ and pick N large enough such that $\sum_{k=N+1}^{\infty} |\frac{x^k}{k!}| < \frac{\varepsilon}{4}$. We can do this as the series for $\exp(x)$ converges absolutely. Now pick $n_0 > N$ large enough so that for each $k = 1, \dots, N$ and each $n \geq n_0$

$$\left| \binom{n}{k} \left(\frac{x}{n}\right)^k - \frac{x^k}{k!} \right| < \frac{\varepsilon}{2N}.$$

We can do this for each $k = 1, \dots, N$ separately and just take the largest n_0 that is needed (N here is fixed, so there are only finitely many k 's). Then for $n \geq n_0$,

$$\begin{aligned} \left| \left(1 + \frac{x}{n}\right)^n - \exp(x) \right| &\leq \sum_{k=1}^N \left| \binom{n}{k} \left(\frac{x}{n}\right)^k - \frac{x^k}{k!} \right| + \sum_{k=N+1}^n \binom{n}{k} \left|\frac{x}{n}\right|^k + \sum_{k=N+1}^{\infty} \left|\frac{x^k}{k!}\right| \\ &\leq \sum_{k=1}^N \left| \binom{n}{k} \left(\frac{x}{n}\right)^k - \frac{x^k}{k!} \right| + \sum_{k=N+1}^n \left|\frac{x^k}{k!}\right| + \sum_{k=N+1}^{\infty} \left|\frac{x^k}{k!}\right| \\ &\leq \sum_{k=1}^N \frac{\varepsilon}{2N} + \frac{\varepsilon}{4} + \frac{\varepsilon}{4} < \varepsilon. \end{aligned}$$

As ε was arbitrary, $(1 + \frac{x}{n})^n \rightarrow \exp(x)$ as $n \rightarrow \infty$. As $e(x)$ we defined as a limit of a subsequence of this sequence, we see that $e(x) = \exp(x)$ also. \square

Note that we could have used this as an alternative proof of Lemma 2. However it uses results on convergence of power series and the *Binomial Theorem*, whereas the proof of Lemma 2 did not.

Once we have (1) for all complex x and y , it is easy to quickly deduce standard results about both \exp and the trigonometric functions.

For example, we can use the power series definitions of $\cos x$ and $\sin x$ to prove for all complex x ,

$$\begin{aligned}e^{ix} &= \cos x + i \sin x, \\e^{-ix} &= \cos x - i \sin x,\end{aligned}$$

or equivalently *define* (for any complex x)

$$\cos x = \frac{e^{ix} + e^{-ix}}{2} \quad \sin x = \frac{e^{ix} - e^{-ix}}{2i}.$$

Addition formulae follow immediately:

$$\begin{aligned}\cos(x+y) + i \sin(x+y) &= e^{i(x+y)} = e^{ix} e^{iy} = (\cos x + i \sin x)(\cos y + i \sin y) \\&= (\cos x \cos y - \sin x \sin y) + i(\cos x \sin y + \sin x \cos y), \\ \cos(x+y) - i \sin(x+y) &= e^{-i(x+y)} = e^{-ix} e^{-iy} = (\cos x - i \sin x)(\cos y - i \sin y) \\&= (\cos x \cos y - \sin x \sin y) - i(\cos x \sin y + \sin x \cos y).\end{aligned}$$

Solving these simultaneous equations (add and subtract) then gives that for all $x, y \in \mathbb{C}$

$$\begin{aligned}\cos(x+y) &= \cos x \cos y - \sin x \sin y, \\ \sin(x+y) &= \cos x \sin y + \sin x \cos y.\end{aligned}$$

Other trigonometric formulae can be derived in a similar manner, valid for all complex numbers.