

# Geometric Group Theory

Cornelia Druțu

University of Oxford

Part C course HT 2023

## A quotation

**William Thurston:** “Mathematics is not about numbers, equations, computations, or algorithms: it is about understanding.”

# Residually finite groups

## Definition

$G$  is **Hopf** if every onto homomorphism  $f : G \rightarrow G$  is an isomorphism.

## Theorem

*A finitely generated residually finite group is Hopf.*

## Corollary

*If  $F(X) = \langle A \rangle$  and  $|A| = |X| = n < \infty$ , then  $F(X) \simeq F(A)$ . (i.e.  $A$  **freely generates**  $F(X)$  i.e.  $A$  is a **free basis** for  $F(X)$ ).*

## Proof.

A bijection  $X \rightarrow A$  extends to  $X \rightarrow F(A)$  which extends to an onto homomorphism  $F(X) \rightarrow F(A)$ . By Universal Property, we have a second onto homomorphism, hence an onto hom.  $F(X) \rightarrow F(A) \rightarrow F(X)$ . Since  $F(X)$  is Hopf, the latter hom. is an isomorphism, hence all are.  $\square$

## Residually finite groups. Simple groups

### Theorem

*A finitely generated residually finite group is Hopf.*

The assumption **finitely generated** cannot be dropped from the theorem.

### Example

- *Consider  $X, Y$  countable.*
- *There exists  $f : X \rightarrow Y$  onto and not injective.*
- *$f$  extends uniquely to an onto group homomorphism  $F(X) \rightarrow F(Y)$ .*

At the other extreme, we have simple groups.

### Definition

$G$  is **simple** if the only normal subgroups are  $\{1\}$  and  $G$ .

# Simple groups

## Example

$\mathbb{Z}/p\mathbb{Z}$ ,  $A_n$ ,  $A_\infty$ ,  $PSL(n, \mathbb{Q})$ , infinite f.g. due to Higman, Thompson, Olshanskii, Burger-Mozes.

## Theorem

*A finitely presented simple group has solvable word problem.*

## Proof.

Let  $w \in F(S)$ . Since  $G$  is simple, if  $w \neq 1$  in  $G$  then  $G = \langle\langle w \rangle\rangle$  and hence  $\langle\langle \{w\} \cup R \rangle\rangle = F(S)$ .

Two procedures:

- 1 Enumerate  $\langle\langle R \rangle\rangle$ . Check if  $w$  appears.
- 2 Enumerate  $\langle\langle \{w\} \cup R \rangle\rangle$ . Check if every  $s \in S$  appears.



# Graphs and Cayley graphs

A main method of investigation is to endow an infinite group with a geometry compatible with its algebraic structure, i.e. invariant by multiplication. This can easily be done for finitely generated groups via Cayley graphs.

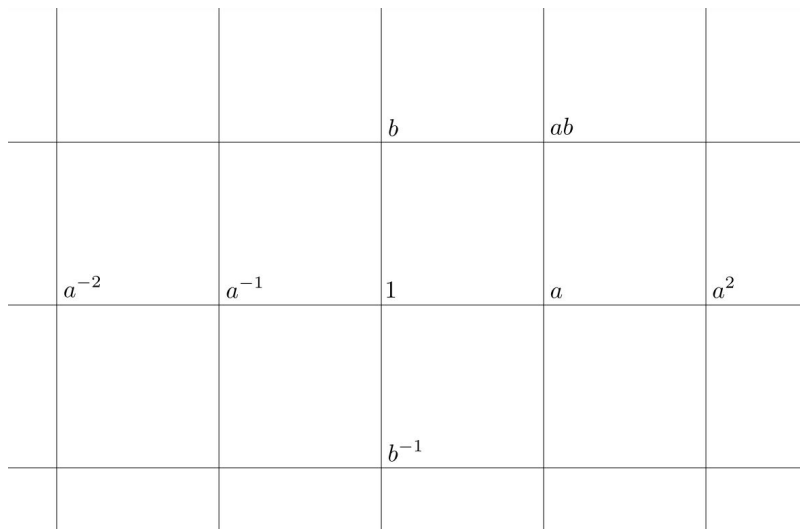
Given a countable group  $G$  and a subset  $S$  such that  $S^{-1} = S$ , the Cayley graph of  $G$  with respect to  $S$ , denoted  $\Gamma(S, G)$ , is a directed/oriented graph with

- set of vertices  $G$ ;
- set of oriented edges  $\{(g, gs) : g \in G, s \in S\}$ ;

We denote an edge  $[g, gs]$ . The underlying non-oriented graph is denoted  $\hat{\Gamma}(S, G)$ .

## Examples of Cayley graphs

1  $\mathbb{Z}^2$  with  $S = \{(\pm 1, 0), (0, \pm 1)\}$



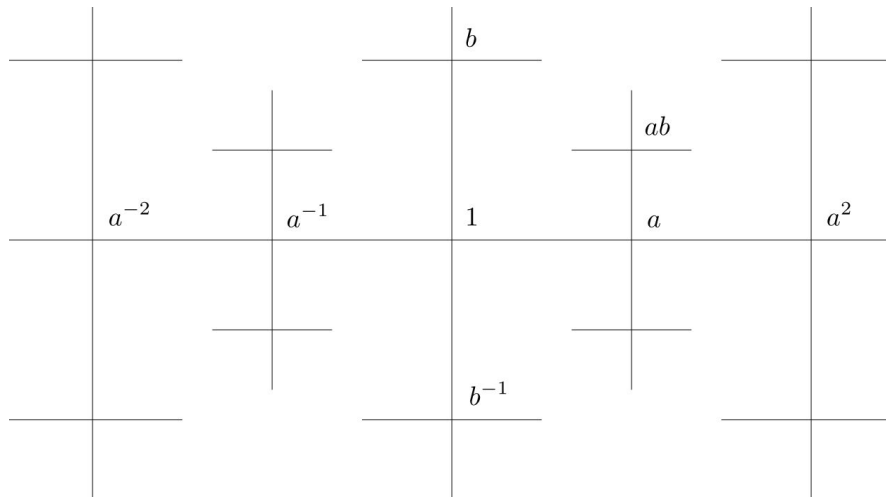
## Examples of Cayley graphs

②  $\mathbb{Z}^2$  with  $S = \{(\pm 1, 0), \pm(1, 1)\}$



## Examples of Cayley graphs

③  $F_2 = F(\{a, b\})$  with  $S = \{a^{\pm 1}, b^{\pm 1}\}$



# Particular features of Cayley graphs

- 1 No **monogons** (edges of the form  $[g, g]$ ) if  $1 \notin S$ .



- 2 No **digons** if, when  $s = s^{-1}$ , we do not list both  $s$  and  $s^{-1}$  in  $S$  (i.e. **no repetitions in  $S$** ).



In other words, this is a **simplicial graph**.

- 3  $\Gamma(S, G)$  is **connected** (i.e. any two vertices can be connected by an edge path) if and only if  $G = \langle S \rangle$ .
- 4
  - a  $\Gamma(S, G)$  is **regular**: the **valency/degree** of every vertex is  $|S|$ .
  - b  $\Gamma(S, G)$  is moreover **locally finite** if and only if  $|S| < \infty$ .

## Particular features of Cayley graphs

- 5 If  $\Gamma(S, G)$  is infinite then it always contains a **bi-infinite geodesic**.



- 6  $\Gamma(S, G)$  always contains a **cycle** (in fact **plenty**) with one exception:  $\Gamma(S, G)$  does not contain a cycle (i.e. it is a **simplicial tree**)  $\iff S = X \sqcup X^{-1}$  and  $G = F(X)$ .

# Cayley Graphs

From now on assume that  $S$  is a finite generating set (with no repetitions),  $1 \notin S$ ,  $S = S^{-1}$ . We endow  $\Gamma(S, G)$  with a metric  $d_S$  as usual:

- each edge has length 1;
- $d_S(x, g)$  is the length of a shortest path from  $x$  to  $g$ .

## Proposition

*The action of  $G$  on its Cayley graph is an action by isometries. The action is free on the vertices. It is free on the whole Cayley graph if and only if no  $s \in S$  is of order 2.*

## Proof.

We have a map

$$G \rightarrow \text{Isom}(\Gamma(S, G)) \quad g \mapsto L_g$$

where  $L_g \in \text{Isom}(\Gamma(S, G))$  extends  $L_g : G \rightarrow G$ ,  $L_g(x) = gx$  to edges.  $\square$

# Cayley Graphs

## Definition

The restriction of  $d_S$  to  $G \times G$  is called the **word metric**.

## Exercises

- $|g|_S := d_S(1, g)$  is *the minimum length of a word  $w$  in  $S$  such that  $g =_G w$* .
- $d_S(g, h)$  is *the minimum length of a word  $w$  in  $S$  such that  $gw =_G h$* .

## Proposition

If  $G = \langle S \rangle = \langle \bar{S} \rangle$  then  $d_S$  and  $d_{\bar{S}}$  are bi-Lipschitz equivalent. That is, there exists  $L > 0$  such that

$$\frac{1}{L}d_S(g, h) \leq d_{\bar{S}}(g, h) \leq Ld_S(g, h)$$

for every  $g, h \in G$ .

# Cayley Graphs

A simplicial tree is a connected graph with no monogons, digons or cycles.

## Theorem

$\hat{\Gamma}(S, G)$  a simplicial tree on which  $G$  acts freely  $\iff S = X \sqcup X^{-1}$ ,  
 $G = F(X)$ .

## Actions on simplicial trees

### Theorem

$\hat{\Gamma}(S, G)$  a simplicial tree on which  $G$  acts freely  $\iff S = X \sqcup X^{-1}$ ,  
 $G = F(X)$ .

### Proof.

( $\Leftarrow$ ) : A cycle would correspond to a reduced word  $w = 1$  in  $F(X)$ .  $\square$